

Design of Secure TRNGs for Cryptography – Past, Present, and Future

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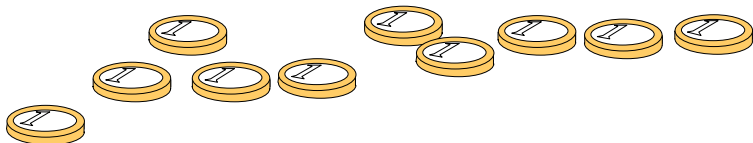
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Random Numbers in Cryptography

- ▶ **(True) Random Number Generator (RNG or TRNG)**
Physical function generating a sequence of random bits or symbols (e.g. groups of bits = numbers)
- ▶ **RNG (or RBG, i.e. Random Bit Generator)**
Essential part of cryptographic systems
- ▶ Today's cryptographic systems mostly implemented in **logic devices** (e.g. smart cards)
- ▶ Challenge: find and exploit **analog sources of randomness in digital devices** using a standard technology (avoid a full custom design)

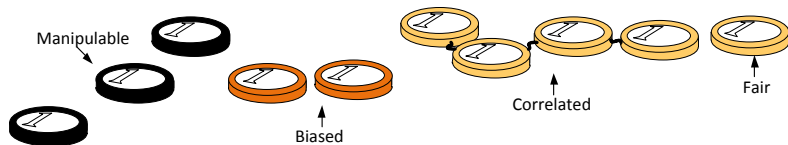
Fair Tossing of Fair Coins



- ▶ Mathematical approach:
 - Considered as an **ideal TRNG**
 - Ten fair coins give **entropy rate of ten bits per trial**
- ▶ Physical approach:
 - What (physically) means **‘fair tossing’**¹ and **‘fair coins’**?
 - What can be the **frequency of trials**?

¹ In fact, mechanical systems are perfectly predictable. Only initial conditions determine the entropy.

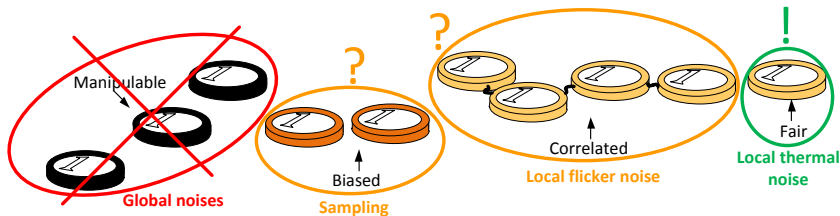
Tossing (Partially) Unfair Coins – Realistic TRNG



- ▶ How much entropy per trial, if:
 - One (independent) fair coin
 - Four correlated coins
 - Two biased coins
 - Three manipulable coins
- ▶ Can the output be manipulable, if the ten coins' values are bit-wise XORed to get **just one output bit**?

Tossing (Partially) Unfair Coins – Realistic TRNG

In the context of oscillator based TRNG:



- ▶ How much entropy per trial, if:
 - One (independent) fair coin
 - Four correlated coins
 - Two biased coins
 - Three manipulable coins
- ▶ Can the output be manipulable, if the ten coins' values are bit-wise XORed to get **just one output bit**?

Conclusions Regarding Our Study Case

- ▶ Design of a TRNG is rather a physical than a mathematical project
- ▶ Physical parameters of the sources of randomness must be thoroughly evaluated:
 - Characteristics of each exploited source of randomness
 - Relationship between individual sources of randomness
 - Distribution of output random values (bias)
 - Correlation or even dependence between output values
 - Manipulability
 - Agility (spectrum)

Random Number Generation and Security

- ▶ Two main security requirements on RNGs:
 - R1: Good statistical properties of the output bitstream
 - R2: Output unpredictability
- ▶ Statistical properties can be easily evaluated using general purpose (black box) statistical tests
- ▶ Unpredictability is more difficult to assess
 - In PRNGs guaranteed by the underlying algorithm – it must be computationally difficult to guess future and past random numbers
 - Approved cryptographic algorithm should be used
 - In TRNGs guaranteed by a sufficient entropy rate per generated random number
 - Approved design approach should be used

Classical versus Modern TRNG Design Approach

- ▶ Recall – two main security requirements on TRNGs:
 - R1: Good statistical properties of the output bitstream
 - R2: Output unpredictability
- ▶ Security evaluation – classical approach:
 - Assess both requirements using statistical tests – insufficient
- ▶ Modern (more stringent) ways of assessing security:
 - Evaluate statistical parameters using statistical tests
 - Evaluate entropy using an entropy estimator (stochastic model)
 - Test online the source of entropy using dedicated statistical tests

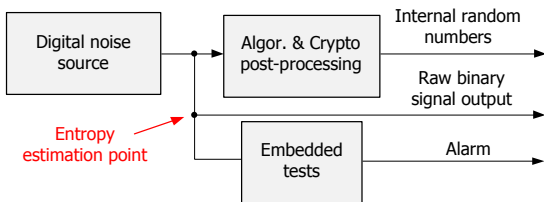
Objectives of the talk

- To discuss modern approaches in the TRNG design
- To illustrate the new methodology on a comprehensive example

Outline

- 1 Sources of randomness in logic devices
- 2 Characterization and quantification of sources of randomness
- 3 From quantification of the source of randomness to dedicated tests
- 4 Conclusions

Contemporary TRNG Design – Recommendations AIS 31



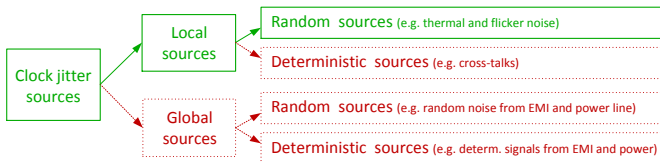
- ▶ Digital noise source
 - Should have as high entropy rate per bit as possible
 - Should enable sufficient bit-rate
 - Shouldn't be manipulable (robustness)
- ▶ Post-processing (optional)
 - Algorithmic – enhances statistics without reducing the entropy
 - Cryptographic – for unpredictability when source of entropy fails
- ▶ Dedicated embedded tests
 - Fast total failure test with low probability of false alarms
 - Online tests detecting quickly and reliably intolerable weaknesses

Sources of Randomness in Logic Devices

- ▶ Commonly used sources related to some physical process, **basically coming from electric noises**
 - **Clock jitter**: short-term variation of an event from its ideal position
 - **Metastability**: ability of an unstable equilibrium electronic state to persist for an indefinite period in a digital system (rare)
 - **Oscillatory metastability**: ability of a bi-stable circuit (e.g. an RS flip-flop) to oscillate for an indefinite period
 - **Initialization of flip-flops**: initialization of a flip-flop (or a memory element) to a random state (after power-up or periodically)
 - **Chaos**: stochastic behavior of a deterministic system which exhibits sensitive dependence on initial conditions

Sources of Randomness: Jittery Clock Signals

- ▶ Clock jitter – the most frequently used in logic devices
- ▶ The jitter in clock generators is caused by ¹
 - Local noise sources
 - Global noise sources



- ▶ **Sources in red are manipulable!**
- ▶ **The entropy must be estimated depending on the local non-manipulable sources (in green)**

¹ B. Valtchanov, A. Aubert, F. Bernard, and V. Fischer, Modeling and observing the jitter in ring oscillators implemented in FPGAs, DDECS 2008

Choice of the Source of Randomness

- ▶ The source of randomness must be **clearly defined, well characterized and quantified**
- ▶ With respect to the entropy harvesting method, it should serve as an **input parameter of the stochastic model**
- ▶ Problem #1: False entropy source
E.g. while claiming to use metastability, the designer uses some other, uncharacterized source of entropy (electric noises)
- ▶ Problem #2: **Entropy overestimation**
The effect of manipulable sources is not excluded from entropy estimation – the general purpose statistical tests are not able to exclude them!

Digitization of the Noise Signal

▶ **Explicite**

- Sampling of a noisy signal
- Counting of random events
- Time-to-digital conversion

▶ **Hidden** (or implicite)

- Conversion of analog electric noises to the timing jitter of the clock signal

▶ Sometimes it is difficult or even **impossible to separate** digitization from the post-processing

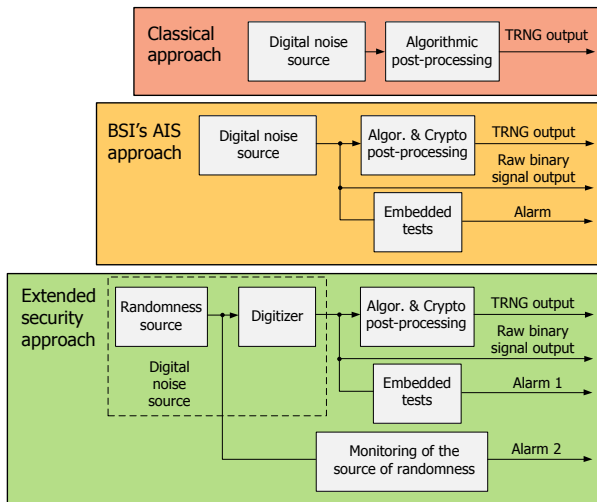
▶ If the digitization is hidden or if it is mixed with the post-processing, the **raw random signal – difficult to determine**

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Secure TRNG Design – Evolution

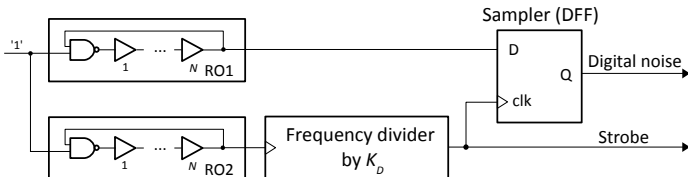
- ▶ TRNG designs should continue to evolve towards security:



Characterization and Quantification of Noise Sources

- ▶ All the sources (and only the sources) that determine the entropy rate at generator's output need to be characterized and quantified
- ▶ Consequently, the noise sources should be characterized and quantified with respect to the stochastic model, which determines the entropy rate
- ▶ Next, we will illustrate this approach on a comprehensive example using an elementary oscillator-based TRNG ...

Elementary Oscillator-Based TRNG (ELO TRNG)

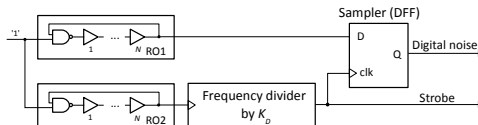


- ▶ First proposed by Fairfield *et al.* ¹
- ▶ Modeled by Baudet *et al.* ² – the entropy depends on the clock jitter coming from the **thermal noise** and the frequencies of the two clock signals
- ▶ The frequency divider determines the sampling period
- ▶ Depending on the jitter size, the K_D value can be very big (greater than 300 000)

¹ R.C. Fairfield, R.L. Mortenson, and K.B. Coulthart. An LSI random number generator (RNG). Advances in Cryptology, 1985

² M. Baudet, D. Lubicz, J. Micolod, and A. Tassiaux. On the security of oscillator-based random number generators. Journal of Cryptology, 2011

ELO TRNG – Security Analysis



- ▶ The effect of the global jitter sources (often neglected!) is significantly reduced by the principle – two identical oscillators are impacted in the same way by the global perturbation signals
- ▶ According to the model, the *lower bound of the Shanon entropy rate* per bit at the generator output is given as:

$$H_{min} \approx 1 - \frac{4}{\pi^2 \ln(2)} e^{-4\pi^2 Q} = 1 - \frac{4}{\pi^2 \ln(2)} e^{-\frac{-4\pi^2 \sigma_{jit}^2 T_2}{T_1^3}} \quad (1)$$

The lower entropy bound is determined by measurable parameters!

- Mean frequencies of the two ring oscillators – T_1, T_2
- Variance of the jitter coming from the **thermal noise** – σ_{jit}^2

Measurement of the Non-Manipulable Clock Jitter ^{1/2}

Algorithm for computing variance V of the jitter¹

- ▶ **Input:** The output sequence $[b_1, \dots, b_n]$ of an elementary TRNG with $K_D = 1$, K , M and N integers ²,
- ▶ **Output:** $V_0 = 4V/T_1^2$ where V is the variance of the jitter accumulated during MT_2 .

Algorithm 1

for $i = 0, \dots, K$ **do**

$S_i \leftarrow [Ni + 1, \dots, Ni + N];$

$c[i] = \mathbb{P}_{S_i}(b_j \neq b_{j+M});$

end for;

$V_0 \leftarrow \frac{1}{K} \sum_{i=0}^K c[i]^2 - \left(\frac{1}{K} \sum_{i=0}^K c[i] \right)^2;$

return: $V_0;$

¹ V. Fischer and D. Lubicz. Embedded evaluation of randomness in oscillator based elementary TRNG. CHES 2014

² In practice, $K \sim 10000$, $N \sim 100$ and $M > N$, we let $M \sim 200 \div 1600$

Measurement of the Non-Manipulable Clock Jitter ^{2/2}

Algorithm 1 – Recall

for $i = 0, \dots, K$ **do**

$S_i \leftarrow [Ni + 1, \dots, Ni + N];$

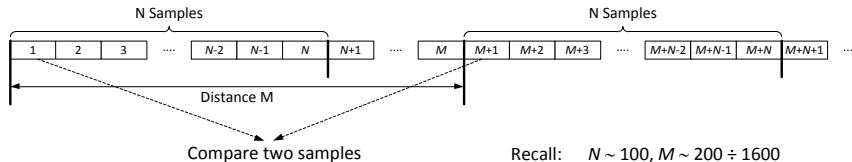
$c[i] = \mathbb{P}_{S_i}(b_j \neq b_{j+M});$

end for;

$V_0 \leftarrow \frac{1}{K} \sum_{i=0}^K c[i]^2 - \left(\frac{1}{K} \sum_{i=0}^K c[i] \right)^2;$

return: $V_0;$

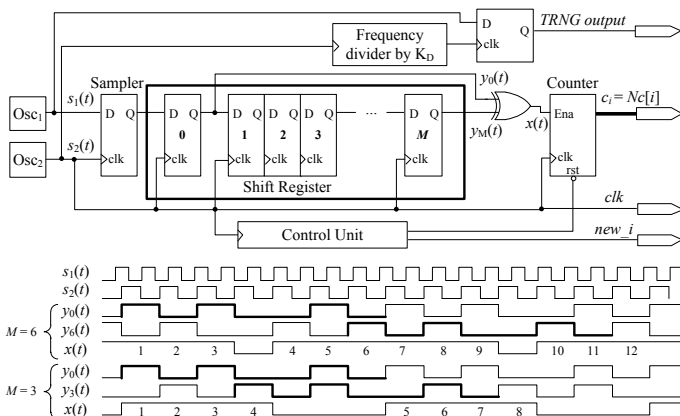
- For all elements from the set S_i compute $c[i] = \frac{\#\{j \in S_{i_0} | b_j \neq b_{j+M}\}}{N}$



Recall: $N \sim 100, M \sim 200 \div 1600$

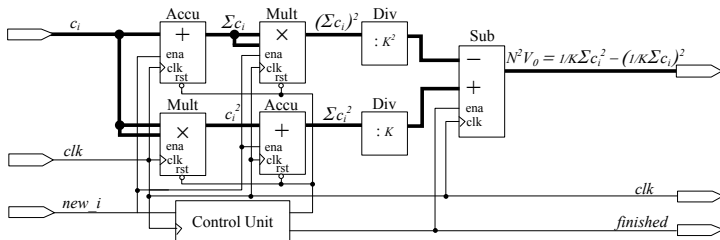
Hardware Implementation of the Jitter Measurement 1/2

- ▶ Jitter measurement circuitry implemented in two blocks
- ▶ The first block computes K successive values $c_i = Nc[i]$



Hardware Implementation of the Jitter Measurement 2/2

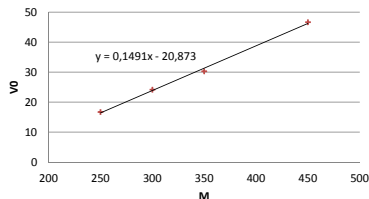
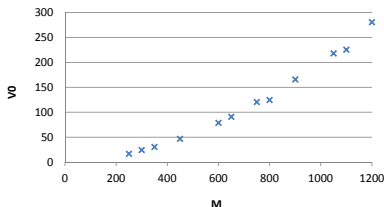
- ▶ Recall: Jitter measurement circuitry implemented in two blocks
- ▶ The second block computes the relative variance $4V/T_1^2$ from K values $c[i]$ according to Algorithm 1



- Summary: Two accumulators, two multipliers, one subtractor, two divisions by shifting right

Evaluation of the Jitter Measurement in Hardware

- ▶ Implementation results in Altera Cyclone III FPGA module
 - The ELO TRNG including jitter measurement circuitry with 32-bit data path occupied:
 - 301 logic cells (LEs),
 - up to 450 memory bits,
 - one DSP block 9x9,
 - four DSP blocks 18x18
- ▶ Jitter measurement results ($250 < M < 1200$, $N \sim 120$ and $K = 8192$)



- From the slope of the measured V_0 for $250 < M < 450$:
Jitter size: $\sigma = 5.01$ ps per period $T_1 = 8.9$ ns.

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Monitoring of the Source of Randomness

- ▶ Monitoring = continuous quantification (embedded measurement) of the noise source
- ▶ The measurement should be performed as close to the source as possible (reduced latency)
- ▶ The impact of the manipulable sources on the measurement results should be avoided
- ▶ The quantified source of randomness should be used
 - As an input for the **stochastic model** for entropy estimation
 - As a basis for the **dedicated stochastic tests** – fast and efficient

Model-Based Entropy Management ^{1/2}

For the previous example:

- ▶ Knowing the size of the jitter, we can now manage the entropy rate at the TRNG output:
 - From Eq. (1), we compute the value of the frequency divider K_D , to ensure that the entropy per bit will always be higher than $H_{min} = 0.997$:

$$K_D > \frac{-\ln\left(\frac{\pi}{2}\sqrt{(1-H_{min})\ln(2)}\right)}{2\pi^2\frac{T_2}{T_1}\frac{\sigma^2}{T_1^2}}$$

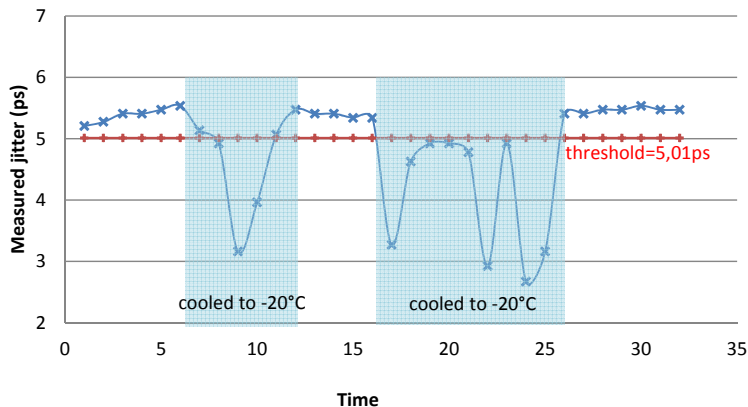
- ▶ For $T_1 = 8.9$ ns, $T_2 = 8.7$ ns, $\sigma = 5.01$ ps and $H_{min} = 0.997$, we get $K_D \approx 430\,000$

Model-Based Entropy Management ^{2/2}

- ▶ **The jitter measurement circuitry can serve for online testing:** for the given K_D , the jitter size σ_c should not drop below 5.01 ps, in order to **guarantee sufficient entropy rate at TRNG output**
- ▶ The proposed dedicated test needs $N \cdot K = 128 \cdot 8192 \approx 1 \cdot 10^6$ periods T_2 to be finished = **less than 3 TRNG output bits!**
- ▶ We observed that the proposed embedded test is **much more conservative** than the tests FIPS 140-1 – the TRNG output passed these tests (and even the tests NIST SP 800-22) for $K_D > 100\,000$
- ▶ **It is sufficient to put a 3-element shift register at the TRNG output, in order to get each output bit continuously tested**

Evaluation of the Method by Attacks

- ▶ Studied attack – jitter reduction by decreasing the temperature
 - The temperature was rapidly changed to -20°C and left to rise back to 21°C for several times.



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Conclusion – TRNGs Suitable for Source Monitoring

- ▶ To comply with the proposed principle of randomness monitoring, the TRNGs must fulfill the following conditions:
 - Their stochastic model must be feasible
 - The model must depend on measurable inputs
- ▶ Not all TRNGs comply with this principle, but many of them do, e.g.:
 - Generators with uniformly distributed clock phases ¹
 - TRNGs with periodically occurring clock phases (coherent sampling)^{2 3}
 - Generators with a transitional oscillatory state ⁴

¹ A. Cherkaoui, V. Fischer, L. Fesquet, A. Aubert: A Very High Speed True Random Number Generator with Entropy Assessment, CHES 2013

² P. Kohlbrenner, K. Gaj: An Embedded True Random Number Generator for FPGAs, ACM/SIGDA FPGA, 2004

³ V. Fischer and M. Drutarovsky: True Random Number Generator Embedded in Reconfigurable Hardware, CHES 2002

⁴ M. Varchola, M. Drutarovsky: New High Entropy Element for FPGA Based True Random Number Generators, CHES 2010

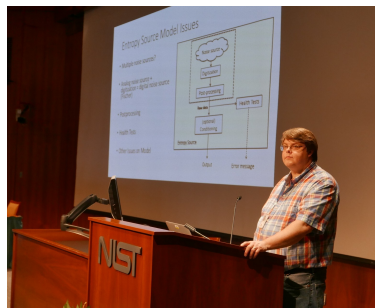
Conclusions

- ▶ We demonstrated that in conjunction with a suitable statistical model, the quantified noise source can be **used to estimate entropy** at the output of the generator
- ▶ We also showed that this entropy estimator can be used to build a **rapid dedicated on-line statistical test** that is perfectly adapted to the generator's principle
- ▶ This approach ensures **high level of security** by rapidly detecting all deviations from the expected behavior

During the NIST RBG workshop in Washington in May 2016

Position of NIST:

- ▶ No set of general-purpose statistical tests can measure the entropy per sample in an arbitrary sequence of values
- ▶ Right way to build a noise source is:
 - Design your noise source
 - Understand it
 - Model it
 - Use your model to estimate entropy
 - Run G-P tests as sanity check
- ▶ We require design documentation and an entropy estimate from designer to support this...
- ▶ ...but we're limited in what resources we can demand for validation testing, and what expertise we can require from labs.



John Kelsey
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HECTOR

Hardware Enabled Crypto and Randomness

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www.hector-project.eu

