

# Exploring different faulting techniques for stressing White-Box cryptography binaries



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WhibOx 2019

- Binary analysis: dynamic fault injection is a powerful way to stress WBC-based solutions,
- Publications in this area remain modest, mostly due to challenging practical realisation,
- Registers, memory access can be changed in runtime leading to exploitable faulty computations,
- Nowadays, a state-of-the-art WBC security analysis must include:
  - static and dynamic fault injections
  - an efficient way to induce dynamic faulty computations: being precise and able to affect large range of instructions
  - a large range of public fault injection attacks exploiting single or multiple faults

**How can this be achieved?**

esDynamic

localhost:8000/lab

ESHARDNewsAndroidDiversReadingTO\_SEEEbookhardeningToolsawesomeLPC1769CryptoLatexNFCVendeursAfterEffectConferencesEBIOS

esDynamic

eshard-notebooks > sva

| Name                             | Last Modified  |
|----------------------------------|----------------|
| 0-SimulatedVulnerabilityAnalysis | 11 minutes ago |
| 1-TracerFrameworks               | 11 minutes ago |
| 2-FaulterFrameworks              | 11 minutes ago |
| 3-SpecialUseCases                | a month ago    |
| HOW                              | 11 minutes ago |
| What is inside.ipynb             | a month ago    |
| SVAModuleQuickStart.pdf          | a month ago    |
| SVANotebooksList.pdf             | a month ago    |

FileEditViewRunKernelTabsSettingsHelp

faulter\_presentation.ipynb

Python 3

## Exploring different faulting techniques for stressing white-box cryptography binaries

```
[1]: from gui_api import *
pv = PrettyView(fault_file='./res/aes_fault_campaign_eax/Log/res.json')
pv.display()
```

### Faults Localisation

PC Register value

Instruction Number

zoom rect

| C source                 | #   | PC     | ASM                      |
|--------------------------|-----|--------|--------------------------|
|                          | 585 | 4009ea | jle 4009b6 <gMul+0x1b>   |
| if ((b & 1) == 1)        | 586 | 4009b6 | movzbl -0x18(%rbp),%eax  |
|                          | 587 | 4009ba | and \$0x1,%eax           |
|                          | 588 | 4009bd | test %eax,%eax           |
| p ^= a;                  | 589 | 4009bf | je 4009c8 <gMul+0x2d>    |
|                          | 590 | 4009c1 | movzbl -0x14(%rbp),%eax  |
|                          | 591 | 4009c5 | xor %al,-0x1(%rbp)       |
| hi_bit_set = (a & 0x80); | 592 | 4009c8 | movzbl -0x14(%rbp),%eax  |
|                          | 593 | 4009cc | and \$0xfffffffff80,%eax |
|                          | 594 | 4009cf | mov %al,-0x9(%rbp)       |
| a <<= 1;                 | 595 | 4009d2 | shlb -0x14(%rbp)         |

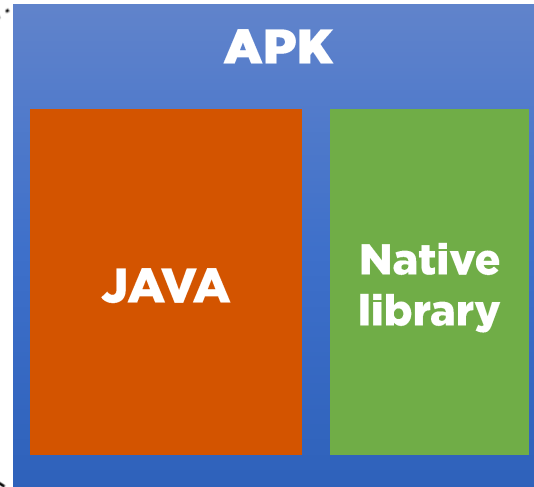
© eshard 2016-2019

# WHERE IS THE WHITE-BOX?

Target the application

Reverse Engineering

White-Box Extraction



Secure functions  
("secure VM")

White-Box  
Cryptography  
Library



Native  
binary file  
(assembly code)



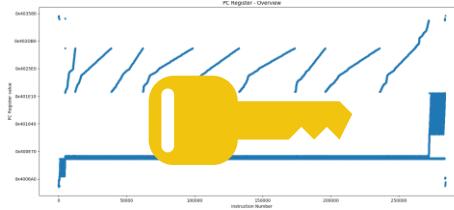
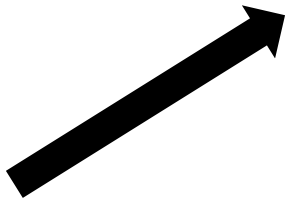
NO SOURCE FILES!

# Why fault a white-box?

# WHY FAULT A WHITE-BOX?



Native  
binary file



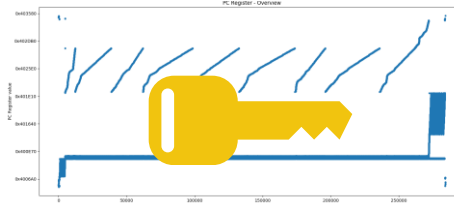
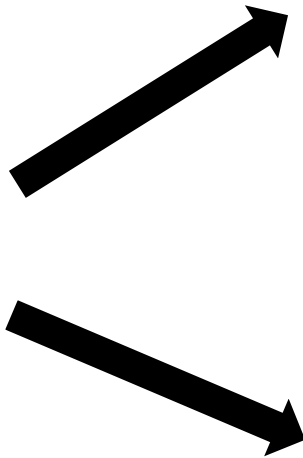
Cryptographic  
attacks

- Differential Fault Analysis (DFA)
- Safe Error
- ...

# WHY FAULT A WHITE-BOX?



Native  
binary file



Cryptographic  
attacks

- Differential Fault Analysis (DFA)
- Safe Error
- ...



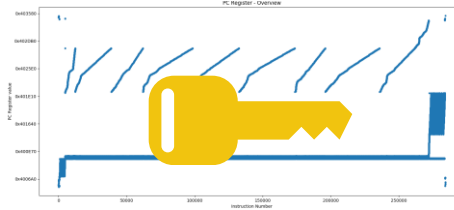
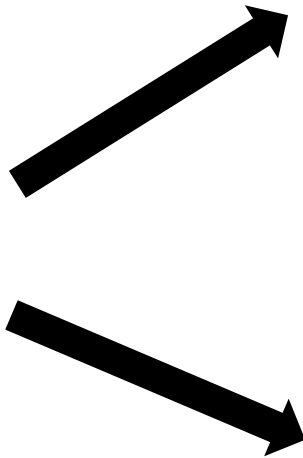
Security  
downgrade

- Defeat integrity mechanisms
- Defeat algorithm protection.  
Stuck mask value to transform a 2<sup>nd</sup> order attack to a 1<sup>st</sup> order
- ...

# WHY FAULT A WHITE-BOX?



Native  
binary file



Cryptographic  
attacks



Security  
downgrade

**How can we find fault  
parameters without  
a COMBINATORIAL  
COMPLEXITY?**



# How to fault a White-Box

```
[3]: slideshow('How to fault a White-Box',  
             './img/presentation_faulter',  
             ['Diapositive%d.PNG' % hit for hit in range(12,21)],  
             slide_width)
```

## ▼ How to fault a White-Box

Slide 0

Slide 1

Slide 2

Slide 3

Slide 4

Slide 5

Slide 6

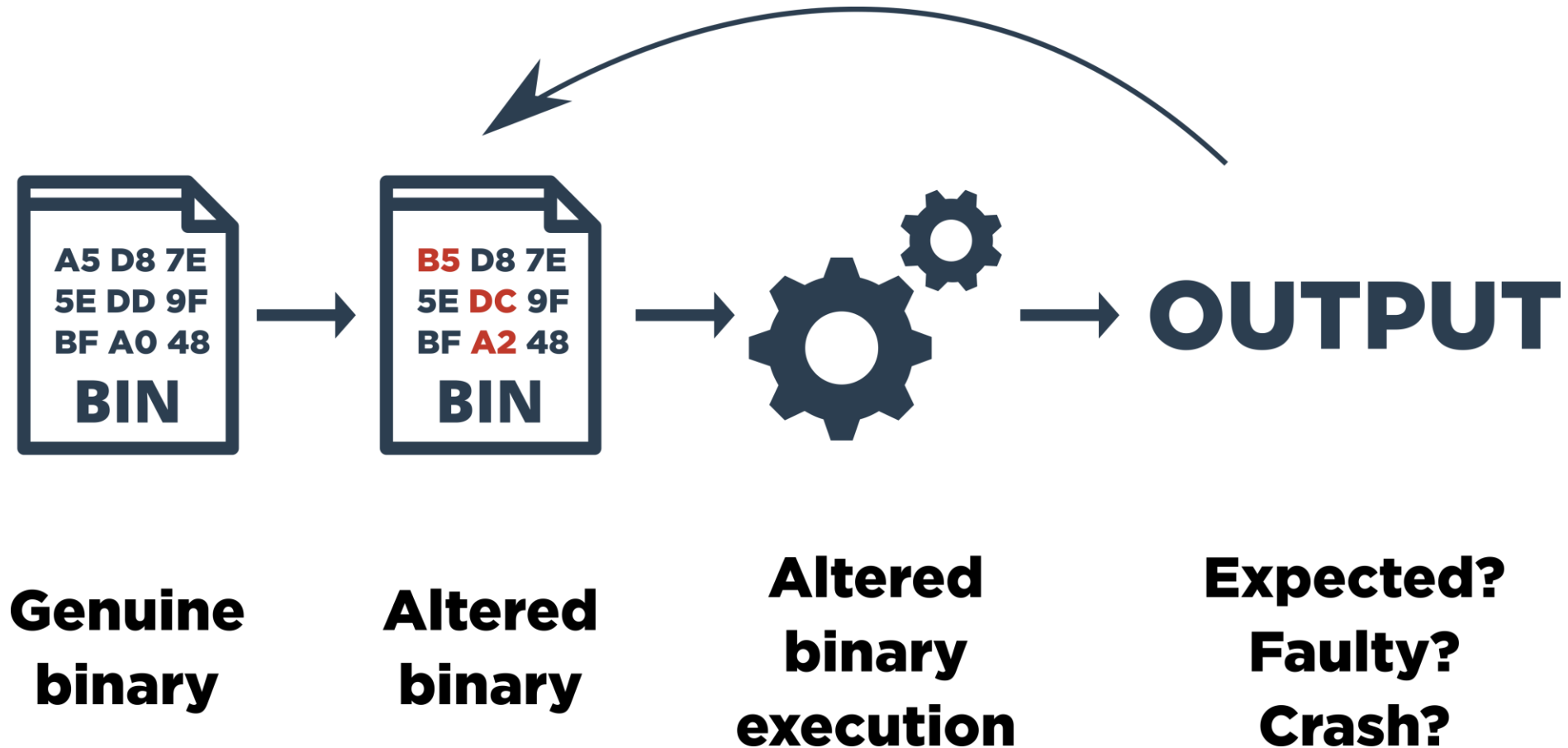
Slide 7

Slide 8

## AGENDA

- Introduction
- **How to fault a White-Box**
- Requirements to fault a White-Box
- Double Fault injection on an AES White-Box
- Conclusion

# STATIC FAULT INJECTION



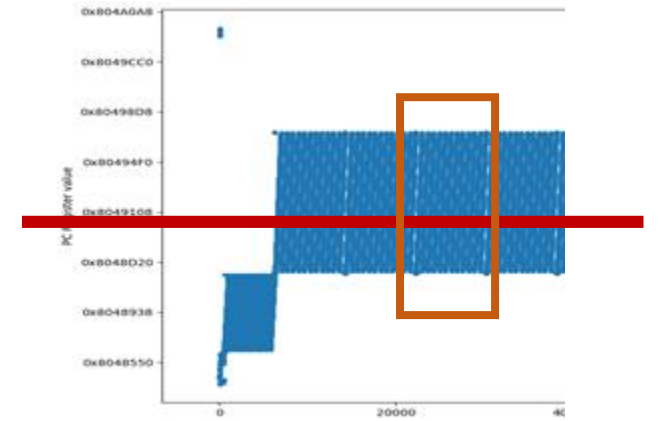
# STATIC FAULT INJECTION

Assets:

- Easy to implement

Drawbacks:

- **Speed**: how to avoid combinatorial complexity with multiple fault injections?
- **Accuracy**: valuable to modify table value, but not disturbing operation execution
- **Anti-Fault countermeasures**: fault easily detected



# STATIC FAULT TOOLS



**Deadpool**  
included in



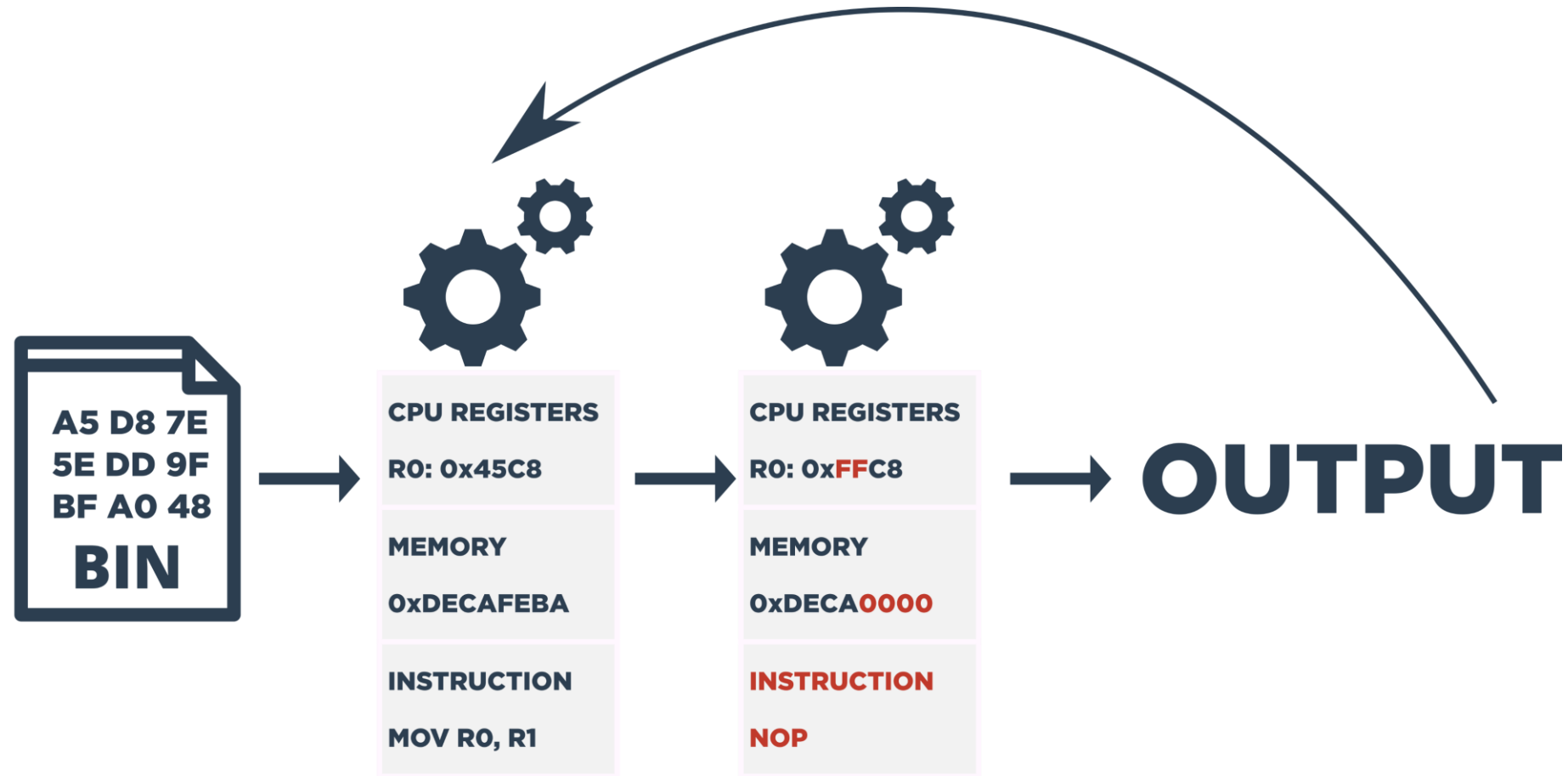
- Python framework
- Tree strategy to inject the faults



**Side Channel Marvels  
Framework**

<https://github.com/SideChannelMarvels>

# DYNAMIC FAULT INJECTION



**Genuine  
binary**

**Binary  
execution**

**Execution  
Alteration**

**Expected?  
Faulty?  
Crash?**

# DYNAMIC FAULT INJECTION

## Assets:

- Accuracy:
  - Alter registers, memory or instructions
  - Multiple fault injection to defeat security countermeasures

## Drawbacks:

- **Fault Model**: which fault effects must be implemented
- **Speed**: how to avoid combinatorial complexity with multiple fault injections?

# DYNAMIC FAULT INJECTION

Assets:

- Open source
- Use the powerful Unicorn Engine...



**Unicorn**

# DYNAMIC FAULT INJECTION

```
# Patch printf and putchar .plt --> return
mu.mem_write(0x80483BC, "\xc3")
mu.mem_write(0x80483EC, "\xc3")
```

```
def hook_mem_access_fault(uc, access, address, size, value, user_data):
    global output, evtId, fault
    evtId += 1
    pc = uc.reg_read(UC_X86_REG_EIP)

    targetId = user_data[0]
    if access == UC_MEM_READ:
        value = u32(uc.mem_read(address, size))
        if should_fault(evtId, targetId, fault, address, size):
            print "FAULTING AT ", targetId
            # Already faulted this time
            fault = False
            # Random bit in this event
            bitfault = 1 << random.randint(0, size*8 -1)
            uc.mem_write(address, pack(value ^ bitfault, size))

    # At this PC the pushes a byte of output to the stack
    if pc == 0x08049e4f:
        output.append(value)
```

Call to external libraries  
must be implemented/patched

Know where to recover the  
ciphertext once the fault was  
injected



# DYNAMIC FAULT INJECTION



**Unicorn**

Assets:

- Open source
- Use the powerful Unicorn Engine...

Drawbacks:

- ... that needs reverse engineering
- Executable/Library must be instrumented by a script
- The Unicorn emulation is slow

# Requirements to fault a White-Box

```
[4]: slideshow('Requirements to fault a White-Box',  
             './img/presentation_faulter',  
             ['Diapositive%d.PNG' % hit for hit in range(21,30)],  
             slide_width)
```

## ▼ Requirements to fault a White-Box

Slide 0

Slide 1

Slide 2

Slide 3

Slide 4

Slide 5

Slide 6

Slide 7

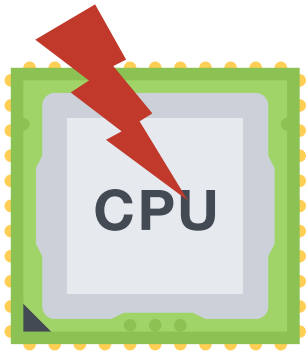
Slide 8

## AGENDA

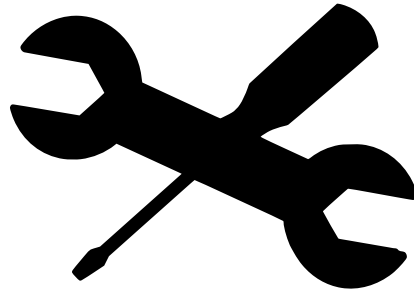
- Introduction
- How to fault a White-Box
- **Requirements to fault a White-Box**
- Double Fault injection on an AES White-Box
- Conclusion

# REQUIREMENTS TO FAULT

Dynamic fault injection



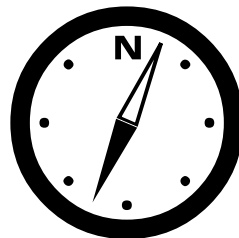
Relevant  
Fault Models



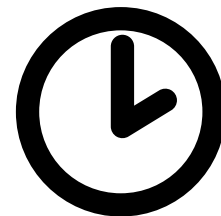
Configuration



Speed



Where?



When?

# REQUIREMENTS TO FAULT

Faults models:

- Register modification

**SET TO ZERO**

**ADD USER VALUE**

**XOR USER VALUE**

**SET USER VALUE**

**SET RANDOM**

# WHAT ARE WE LOOKING FOR?

Faults models:

- Register modification
- Data Flow Modification

**SET TO ZERO**

**ADD USER VALUE**

**XOR USER VALUE**

**SET USER VALUE**

**SET RANDOM**

# REQUIREMENTS TO FAULT

Faults models:

- Register modification
  - Data Flow Modification
  - Control Flow Modification with Program Counter Register

**SET TO ZERO**

**ADD USER VALUE**

**XOR USER VALUE**

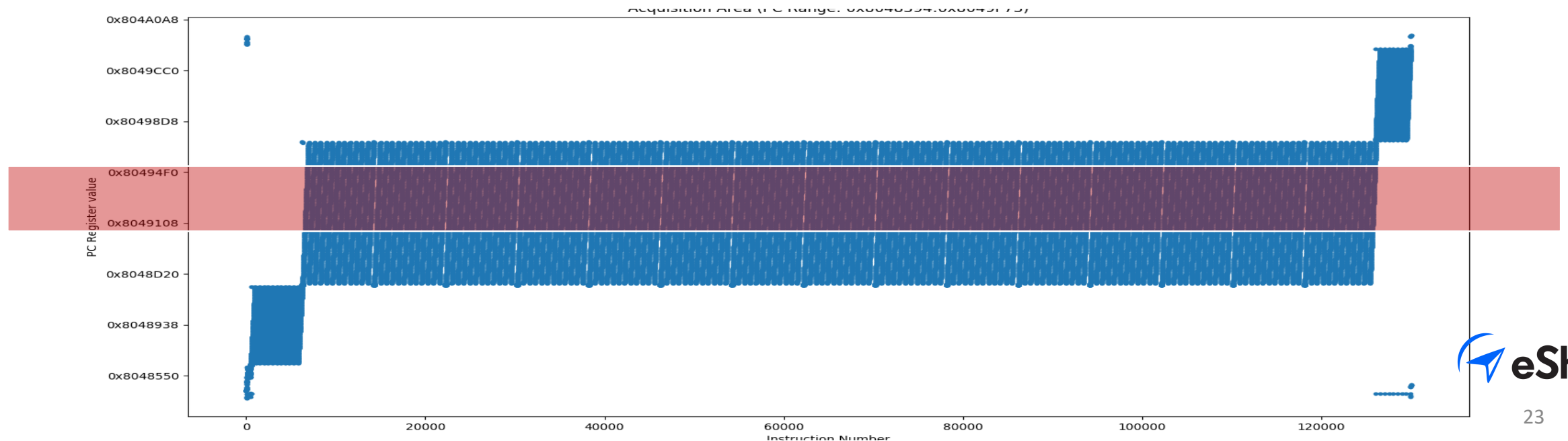
**SET USER VALUE**

**SET RANDOM**

# REQUIREMENTS TO FAULT

Where to fault:

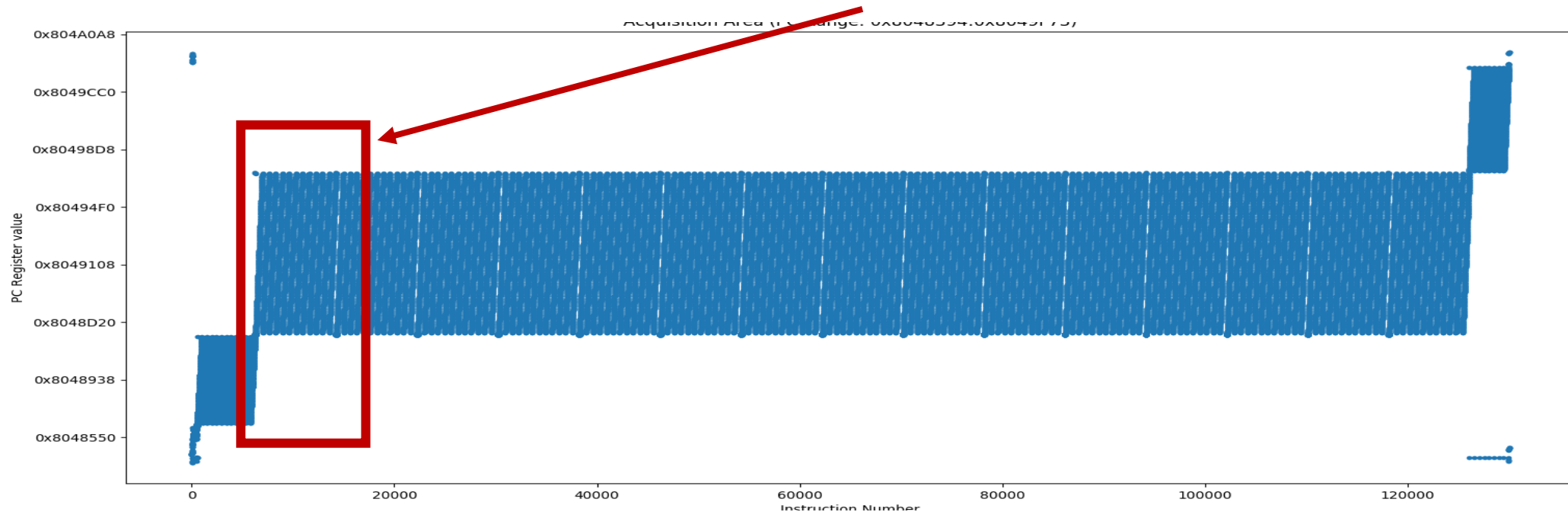
- Filtering based on:
  - Program Counter
  - Kind of instruction: mov, add ...



# REQUIREMENTS TO FAULT

When to fault: pattern detector

- **We must start the faults when the correct PCs are reached ...**
- **... and we also stop them in the same way**

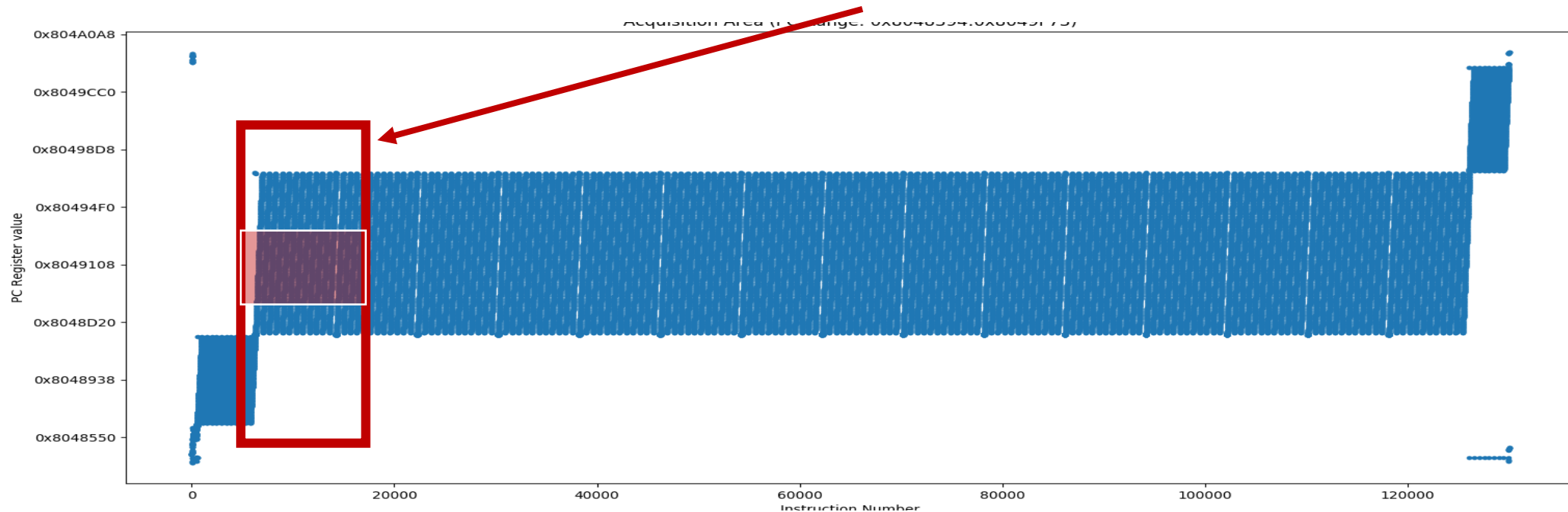




# REQUIREMENTS TO FAULT

When to fault: pattern detector

- **It can be combined with a filtering based on Program Counter value or Instruction kind**



# REQUIREMENTS TO FAULT

 **eShard**  
**esFaulter**

included in

**esDynamic  
Platform**



- Dynamic register modification
  - Data flow disturbance
  - Control flow disturbance
- Multi-fault injection
- Fault&trace capabilities
- Filtering and trig&act capabilities



 **QEMU**

- X86, x86\_64, ARM support
- Takes advantage of Qemu **speed**

# Double Fault injection on an AES White-Box

```
[5]: slideshow('Introduction',  
             './img/presentation_faulter',  
             ['Diapositive%d.PNG' % hit for hit in range(30,34)],  
             slide_width)
```

## ▼ Introduction

Slide 0

Slide 1

Slide 2

Slide 3

## AGENDA

- Introduction
- How to fault a White-Box
- Requirements to fault a White-Box
- **Double Fault injection on an AES White-Box**
- Conclusion

# THE CHALLENGE

- Attack an AES White-Box implementation
- Configuration:
  - CPU i7-7560U, 2.4GHz dual core
  - 16 GB of RAM (we not need so much)
  - SSD NVMe
- Double fault injection
- Key recovery from the faulty outputs with a DFA of Piret (with 4 modified bytes in a specific way)

# THE CHALLENGE

- Attack an AES White-Box implementation
- Double fault injection
- Key recovery from the faulty outputs with a DFA of Piret (with 4 modified bytes in a specific way)

**How can we find faults parameters without a COMBINATORIAL COMPLEXITY ?**

*$nb\_ins \times nb\_fault\_model \times nb\_target \times nb\_input \times nb\_area$*

# ILLUSTRATION

AES-128  
X86-64 architecture

```
$/cipheraes 06 1F C9 F5 88 B2 F9 D2 00 19 86 82 2C 12 11 79  
message: 061fc9f588b2f9d2001986822c121179  
cipher: 14ed01ea7ce2a551c9791ae85c7cecf4
```

Differential Fault Analysis with a double fault injection attack:

- Data flow disturbance
- Control flow disturbance

Our target is a x86-64 architecture executable.

[6]: `!file ./binary/cipheraes`

```
./binary/cipheraes: ELF 64-bit LSB executable, x86-64, version 1 (SYSV), dynamically linked, interpreter /lib64/ld-linux-x86-64.so.2, for GNU/Linux 2.6.32, BuildID[sha1]=b4bdf121200bbae155c26bc03d474e9f29eb980f, stripped, with debug_info
```

Let's see how to use it.

[7]: `!./binary/cipheraes`

```
cipheraes B01 B02 B03 B04 B05 B06 B07 B08 B09 B10 B11 B12 B13 B14 B15 B16, BXX is hexadecimal arguments problem
```

We need to provide a 16-byte hexadecimal input, and then get the ciphered message.

[8]: `!./binary/cipheraes 061FC9F588B2F9D2001986822C121179`

```
cipheraes B01 B02 B03 B04 B05 B06 B07 B08 B09 B10 B11 B12 B13 B14 B15 B16, BXX is hexadecimal arguments problem
```

The input format must be carefully respected in order to get the ciphered message.

[9]: `!./binary/cipheraes 06 1F C9 F5 88 B2 F9 D2 00 19 86 82 2C 12 11 79`

```
message: 061fc9f588b2f9d2001986822c121179
cipher: 14ed01ea7ce2a551c9791ae85c7cecf4
```

We understood how the binary works, let see now how to use esFaulter.

## Where to attack?

```
[10]: slideshow('Where to attack?',
               './img/presentation_faulter',
               ['Diapositive%d.PNG' % hit for hit in range(34,36)],
               slide width)
```

### ▼ Where to attack?

Slide 0

Slide 1

# FIRST STEP

## Where inject faults?

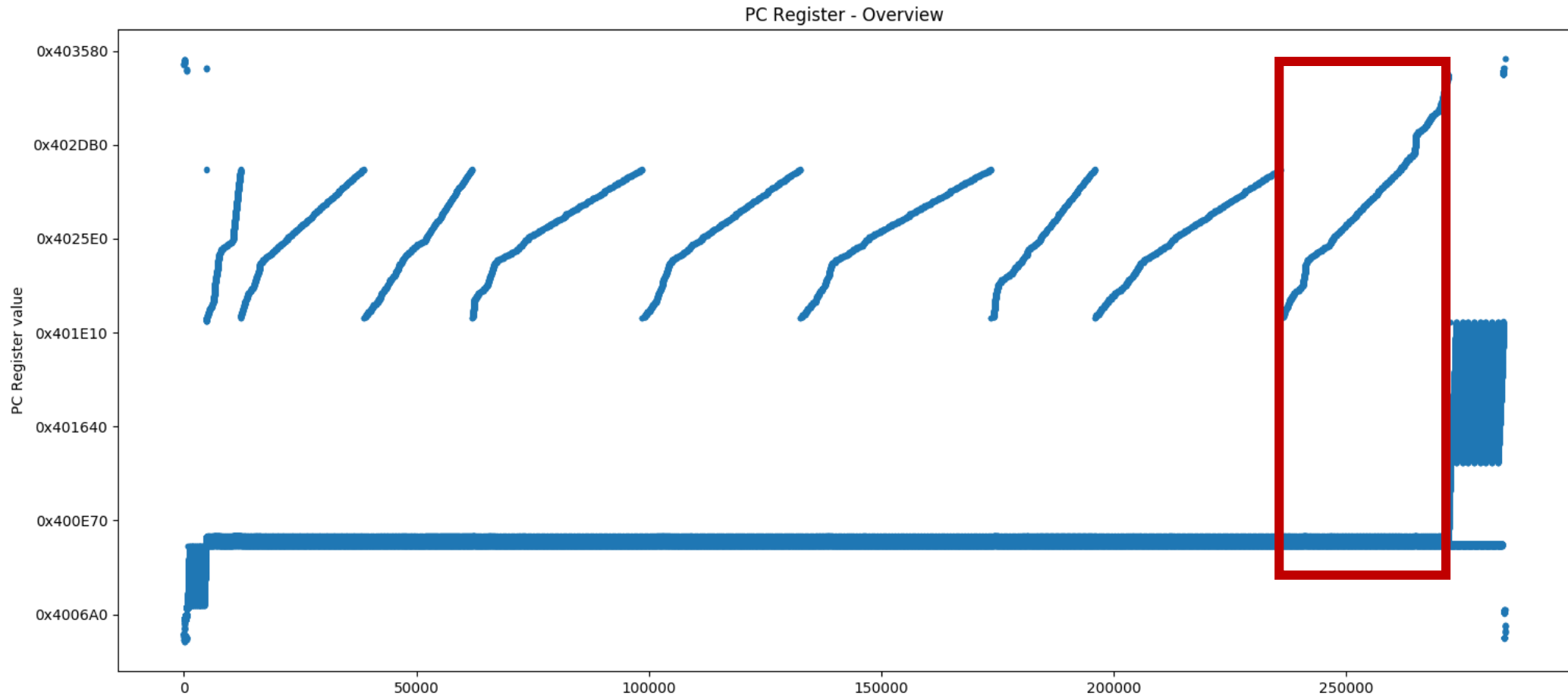
- only the binary itself, not external system libraries

## How to know where to inject faults?

- Trace memory access or registers
- Display them to see distinguishable patterns
- Program Counter (PC), address of executed instruction, tracing is a good start



# SIMPLE FAULT INJECTION



**AES round patterns? Let's attack the last one**

Before performing fault injections, we need to know where to inject them.

We are going to use esTracer to acquire the PC register evolution, and then, we analyze its graphical representation.

First, we configure the binary handler configuration object. It eases the management of the binary to attack.

```
[11]: from essva.binary_handler import BinaryHandlerConfiguration
import re

bin_hand = BinaryHandlerConfiguration()

bin_hand.algo = "AES-128"
bin_hand.cmd = "./binary/cipheraes"
```

We need to indicate how to give a correct input to the program, and as well how to interpret its response.

```
[12]: def process_input(block):
        return " ".join("%02X" % hit for hit in bytes.fromhex(block)).lower()

    def process_output(output):
        return re.match("message: [a-f,0-9]{32}\\ncipher: ([a-f,0-9]{32})?", output).group(1)

    bin_hand.process_input = process_input
    bin_hand.process_output = process_output
```

We are ready to use esTracer!!!

We trace the program scanning only the PC register.

```
[13]: from essva.tracer import Tracer
      tracer = Tracer(bin_hand)
```

By default, it will trace everything, so we configure it:

- we want to trace only the PC register.
- we focus the tracing on the binary itself, the PC range value can be retrieved for instance with readelf.

```
[14]: tracer.registers = ["pc"]
      tracer.filt_pre.in_pc.ranges = [[0x400458, 0x4034DC]]
```

We indicate where the traces will be generated.

```
[15]: tracer.directory_out = "traces"
```

We generate one trace. esTracer can:

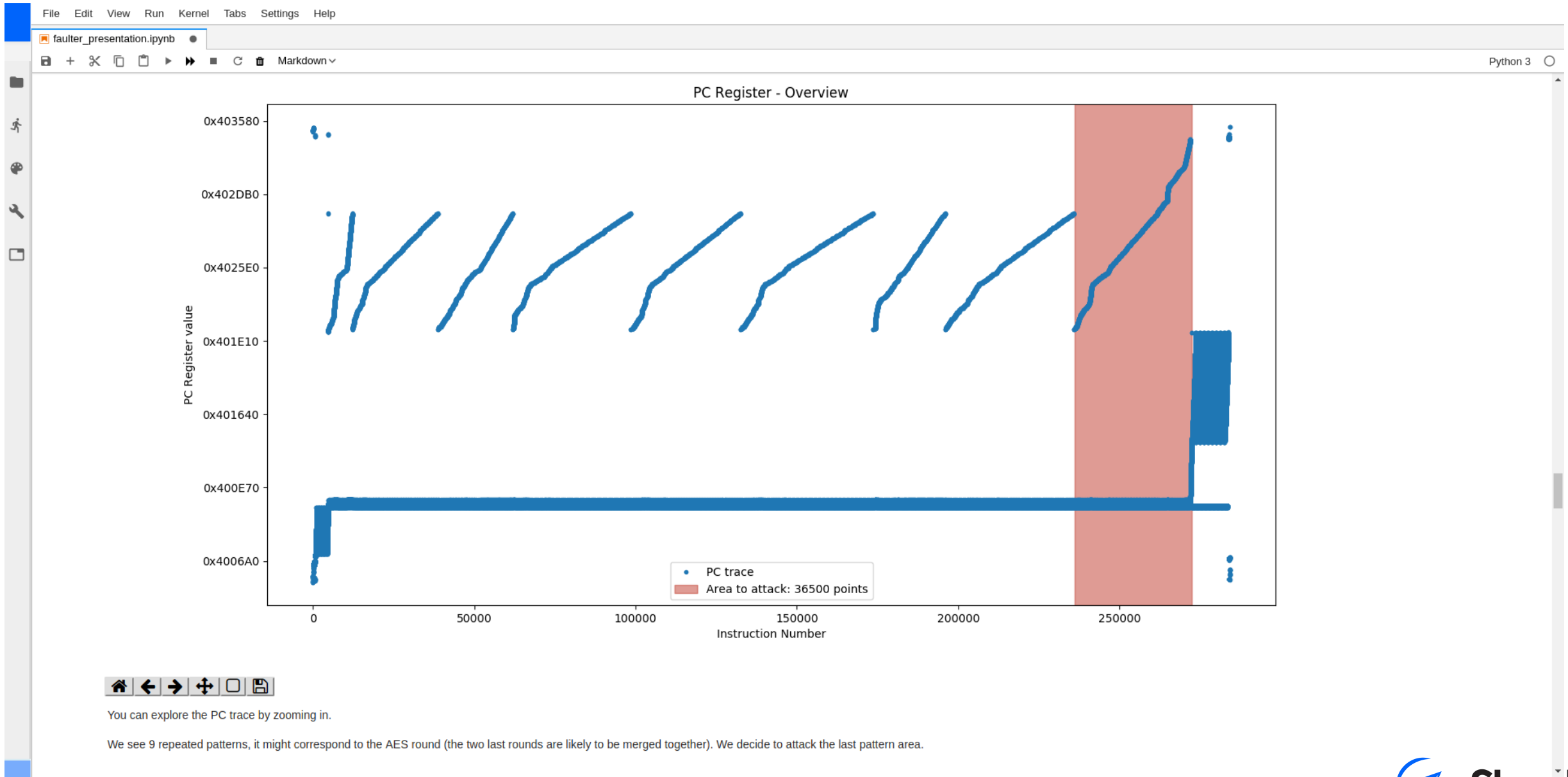
- either automatically generate a random input,
- either use a list of user input. We will use this option.

```
[16]: plain_input = "061FC9F588B2F9D2001986822C121179"
```

```
[17]: tracer.generate_traces([plain_input])
```

Generating 1 trace(s). Please wait...

Processing 100.0 %  Elapsed: 00:00:00 Remaining: 00:00:00



## Fault injection campaign

We will inject a simple fault attack to perform a DFA of Piret, it means we have to obtain outputs with 4 faulty bytes.

We import the esFaulter object, and link it to the binary handler configuration.

```
[21]: from essva.faulter import Faulter
import essva.sva_constants as const

faulter = Faulter(bin_hand)
```

We indicate where to generate the log.

```
[22]: faulter.directory_out = "faulter_outputs_one_fault"
```

We use two trig&act to focus on the last pattern area:

- when the PC 0x400cd2 was called 5297 times, we start to inject faults
- when the PC 0x400c5c was called 4938 times, we stop the fault injection

See the presentation "Optimize your binary tracing: an example with an ECDSA implementation" for more details about the trig&act.

```
[23]: nb_inst_start, nb_inst_end = 236000, 272500
pc_start, pc_end = int(pc_np[nb_inst_start]), int(pc_np[nb_inst_end])
pc_start_count, pc_end_count = len(np.where(pc_np[:nb_inst_start+1] == pc_start)[0]), len(np.where(pc_np[:nb_inst_end+1] == pc_end)

[24]: trig_and_act_start_fault_injection = [pc_start, pc_start_count, const.TAA_F_START_FAULT]
trig_and_act_stop_fault_injection = [pc_end, pc_end_count, const.TAA_F_STOP_FAULT]
```

File Edit View Run Kernel Tabs Settings Help

faulter\_presentation.ipynb

Markdown

We are going to fault all the registers, excepting the PC.

```
[ ]: reg_lst = ["rax", "rcx", "rdx", "rbx", "rsp", "rbp",  
              "rsi", "rdi", "r8", "r9", "r10", "r11", "r12",  
              "r13", "r14", "r15"]
```

To configure a fault, we use the FaultCfgMono object:

- we indicate the register to fault
- we indicate which fault model to use, in our case we set the register to zero
- we indicate when the fault must be injected with the two trig&act created previously

Then, it is linked to esFaulter with `faulter.faults_opt` that is a list of FaultCfgMono object:

- in that way, you can easily perform simple or multiple fault injection
- esFaulter analyzes the configuration of each fault to make **automatic** and **exhaustive** fault injection

```
[ ]: from essva.faulter import FaultCfgMono  
  
for register in reg_lst:  
    fault_1 = FaultCfgMono(bin_hand)  
  
    fault_1.trig_and_acts = [trig_and_act_start_fault_injection,  
                           trig_and_act_stop_fault_injection]  
  
    fault_1.register = register  
    fault_1.model = const.FM_SET_ZERO  
  
    faulter.faults_opt = [fault_1]  
    faulter.directory_out = "fault_1_%s" % register  
  
    faulter.generate_faults(plain_input)
```

```
[28]: slideshow('Simple Fault Injection',  
              './img/presentation_faulter',  
              ['Diapositive%d.PNG' % hit for hit in range(36,39)],  
              slide_width)
```

### ▼ Simple Fault Injection

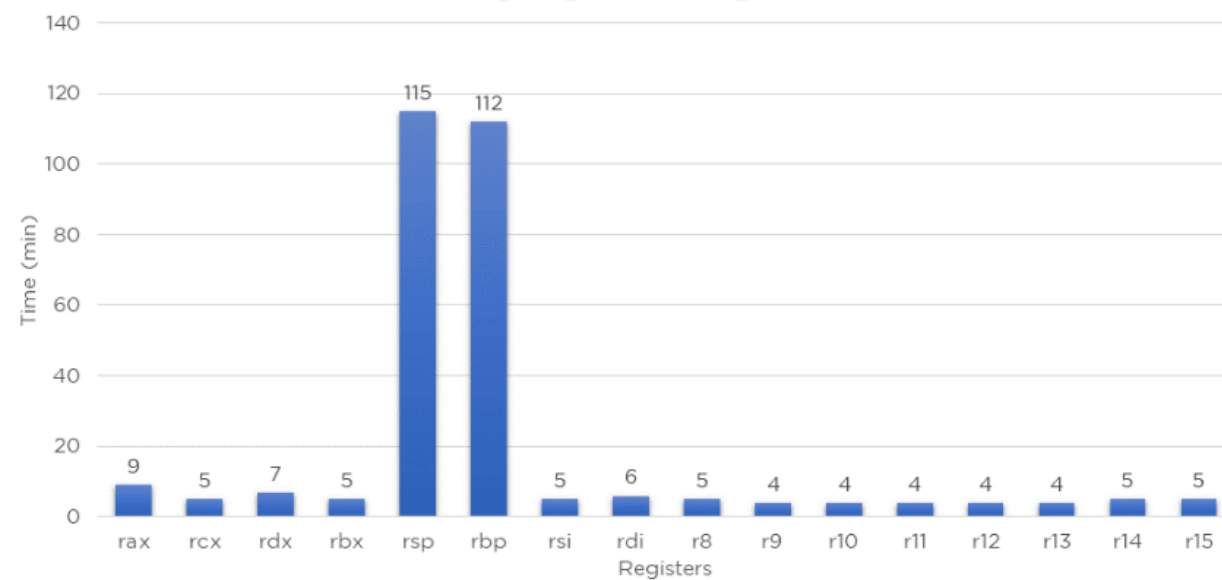
Slide 0

Slide 1

Slide 2

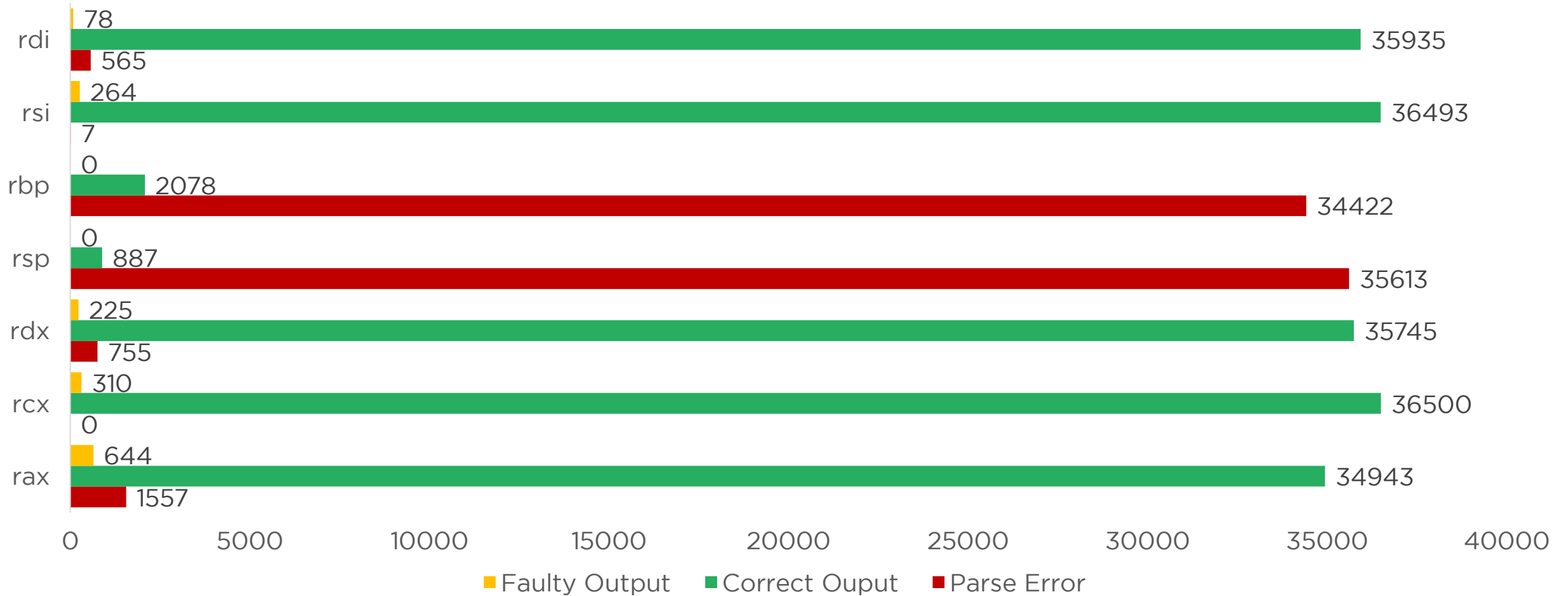
# SIMPLE FAULT INJECTION

## Fault Campaign Timing Overview



# SIMPLE FAULT INJECTION

## Output Classification



No effect for the other registers (rbx, r8, r9, r10, r11, r12, r13, r14, r15)



# SIMPLE FAULT INJECTION



14 campaigns faulting one register (rsp/rbp not included)



511,000 injected faults



~76 min (multi-thread not used)



~112 injected faults by second

## RSP Register Fault analysis

We analyze the log for the rsp fault injection. Instead of looking at the log file, we can use a log parser to have a pretty view.

```
[ ]: from essva.esdynamic.faulter_output import FaulterLogParser
import glob
import re

logfile = glob.glob("./fault_1_rsp/*.bin")[0]
parser = FaulterLogParser(logfile, output_parser = bin_hand.process_output)
parser.parse_file()
# Print faulty output
parser.get_faulty_output()
```

Most of the fault leads to the crash of the application. It is an expected behavior since we set the stack pointer to 0. The same behavior happens for RBP, which is the base pointer.

## RAX Register Fault analysis

We analyze the log for the rax fault injection.

```
[ ]: from essva.esdynamic.faulter_output import FaulterLogParser
import glob
import re

logfile = glob.glob("./fault_1_rax/*.bin")[0]
parser = FaulterLogParser(logfile, output_parser = bin_hand.process_output)
parser.parse_file()
# Print faulty output
parser.get_faulty_output()
```

The program never crashed. But, all the faults have the same effects: we obtain a block of 16 bytes set to 0.

```
[31]: slideshow('Introduction',  
              './img/presentation_faulter',  
              ['Diapositive%d.PNG' % hit for hit in range(39,43)],  
              slide_width)
```

## ▼ Introduction

Slide 0

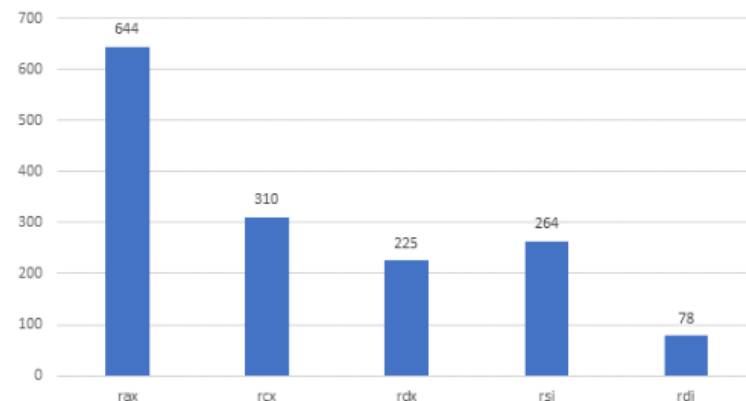
Slide 1

Slide 2

Slide 3

# SIMPLE FAULT INJECTION

Faulty Output



All faulty outputs return  
a ciphertext all bytes  
set to zero!

**How can we interpret this behaviour?**

# HOW TO INTERPRET THE FAULTS?

Correct/  
Faulty output  
analysis



Execute algorithm:

- To recover the key
- Or to detect in which round the fault was injected
- A lot of public algorithm available, but if they fail it gives no information

# HOW TO INTERPRET THE FAULTS?

Correct/  
Faulty output  
analysis



Execute algorithm:

- To recover the key
- Or to detect in which round the fault was injected
- A lot of public algorithm available, but if they fail it gives no information

Reverse  
engineering



- Understand the effect of the fault on the program execution
- Give a way to understand very accurately the fault but it requires reverse engineering skills

# HOW TO INTERPRET THE FAULTS?

Correct/  
Faulty output  
analysis



Execute algorithm:

- To recover the key
- Or to detect in which round the fault was injected
- A lot of public algorithm available, but if they fail it gives no information

Reverse  
engineering



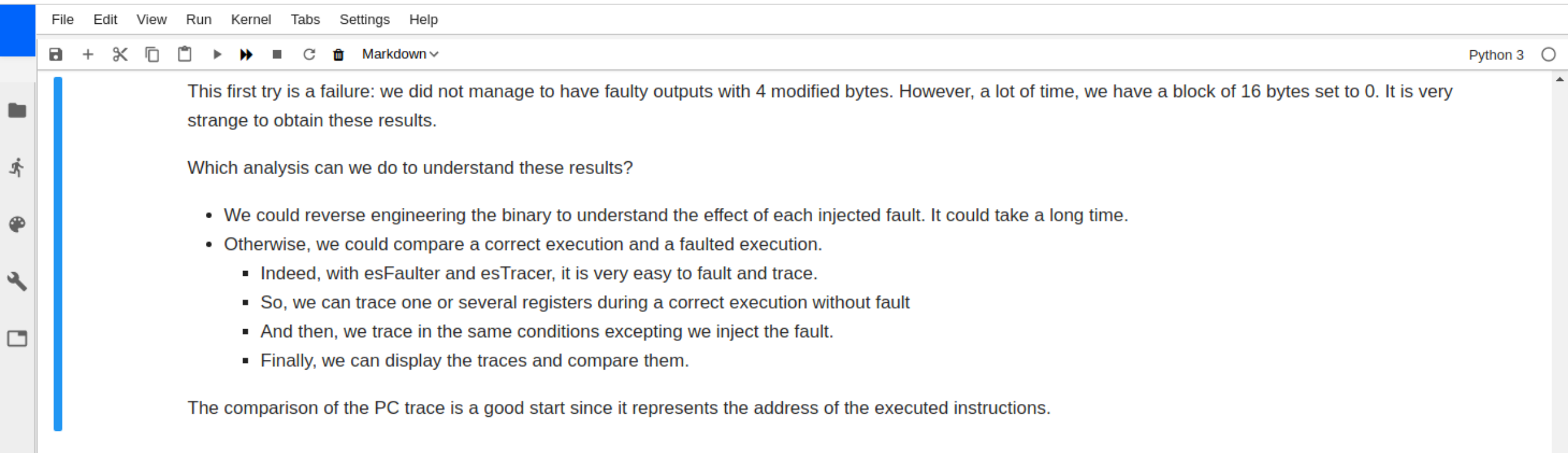
- Understand the effect of the fault on the program execution
- Give a way to understand very accurately the fault but it requires reverse engineering skills

Fault & Trace



- Fault and trace at the same time (memory access, Program Counter registers ...)
- Give a visual way to understand accurately the fault effect without reverse engineering skills

## LET TRY THIS WAY



## Fault & trace

```
[58]: slideshow('Introduction',  
             './img/presentation_faulter',  
             ['Diapositive%d.PNG' % hit for hit in [45, 47, 48]],  
             slide_width)
```

### ▼ Introduction

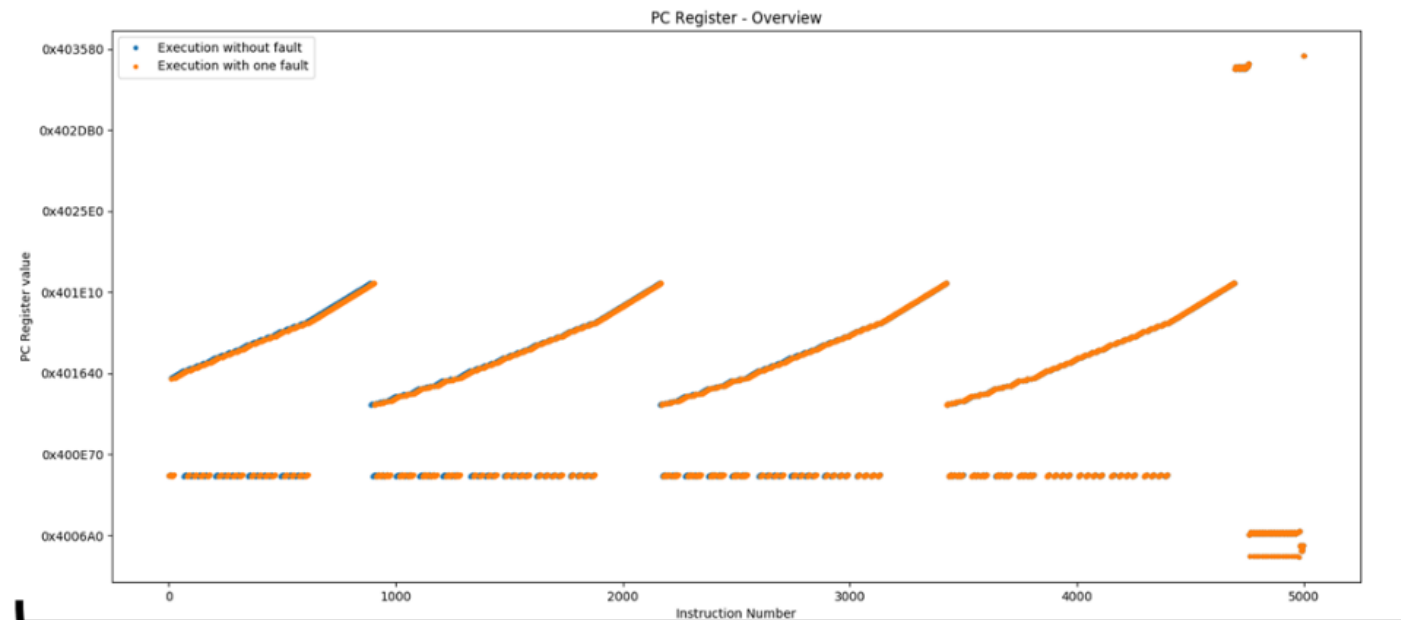
Slide 0

Slide 1

Slide 2

# FAULT & TRACE

Fault and trace the PC register to see the executed instructions



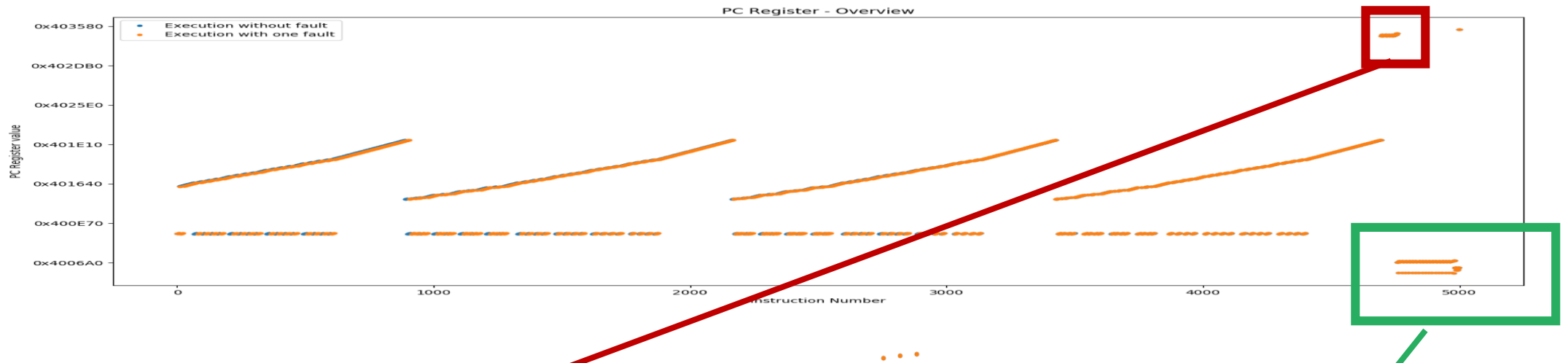
End of the AES execution

45



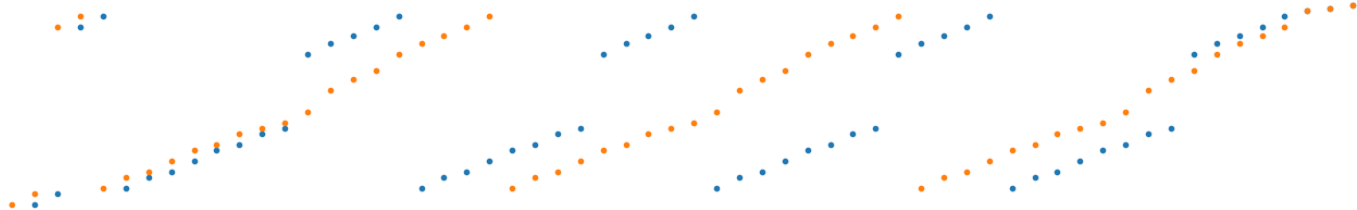
# FAULT & TRACE

Fault and trace the PC register to see the executed instructions



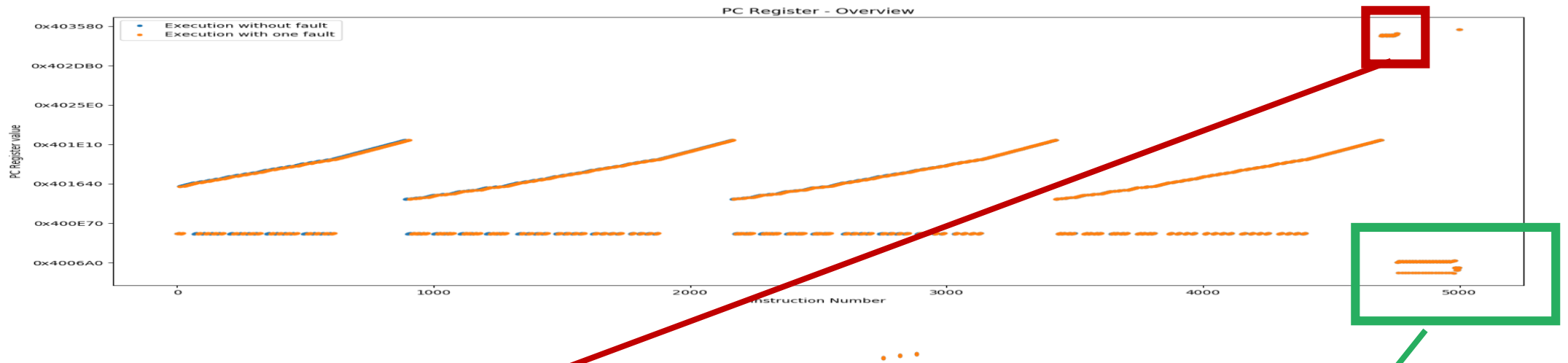
Traces are different

Traces are identical



# FAULT & TRACE

Fault and trace the PC register to see the executed instructions



Traces are different

Traces are identical

**LET SEE WHAT IS  
THE CODE IN THIS AREA**

## Fault & trace

It is very easy to fault and trace:

- you create a tracer object
- you configure it
- and you link it to the faultier
- automatically, for each injected fault, a trace will be generated.

We have already a tracer recording the PC, so we will use it:

```
[32]: faulter.tracer = tracer
```

We analyze the trace when we fault rax register.

- We will not once again inject 36500 faults, and create 36500 traces.
- We only need to inject one fault leading to a zero output.

So, we analyze a log to get a PC value triggering a zero output.

[illegible]

```
[34]: '0x40138c'
```

This PC value might be called several time in the area we attack, so we use a trig&act to fault it just once.

```
[35]: nb_inst_start_f1 = np.where(pc_np == 0x40138c)[0][0]
      nb_inst_end_f1 = nb_inst_start_f1 + 1
      pc_f1_start, pc_f1_end = int(pc_np[nb_inst_start_f1]), int(pc_np[nb_inst_end_f1])
      pc_f1_start_count, pc_f1_end_count = len(np.where(pc_np[:nb_inst_start_f1+1] == pc_f1_start)[0]), len(np.where(pc_np[:nb_inst_end_f1+1] == pc_f1_end)[0])
```

We modify the fault options:

- we set the PC we want to fault
- we set the trig&act

```
[36]: fault_1.filt_pre.in_pc.pts = [pc_to_attack]
      fault_1.trig_and_acts = [[pc_start, pc_start_count, const.TAA_F_START_FAULT],
                             [pc_end, pc_end_count, const.TAA_F_STOP_FAULT]]
      fault_1.register = "rax"
```

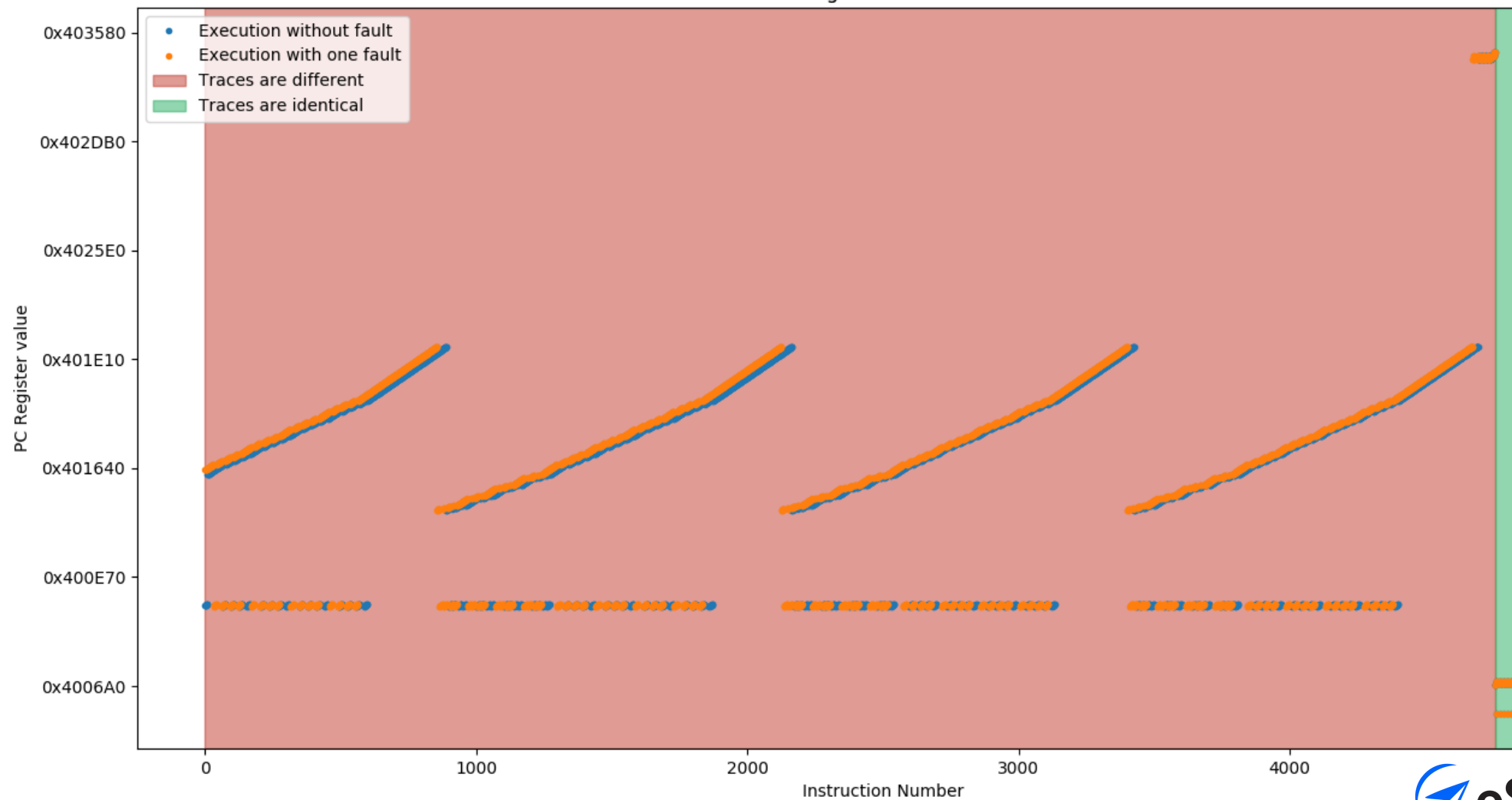
We inject the fault.

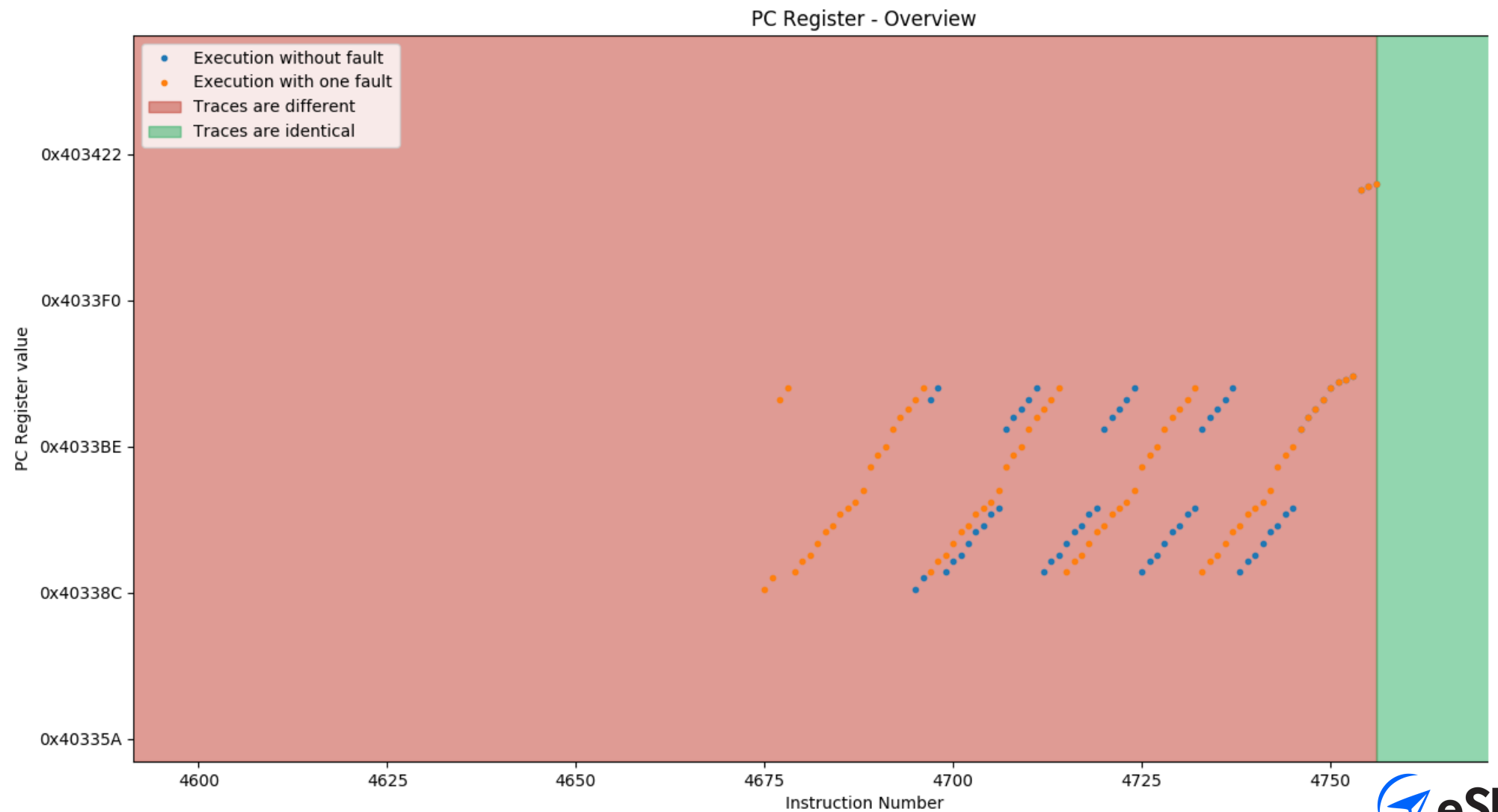
```
[37]: fault_1.directory_out = "fault_1_rax_with_trace"
      fault_1.generate_faults(plain_input)
```

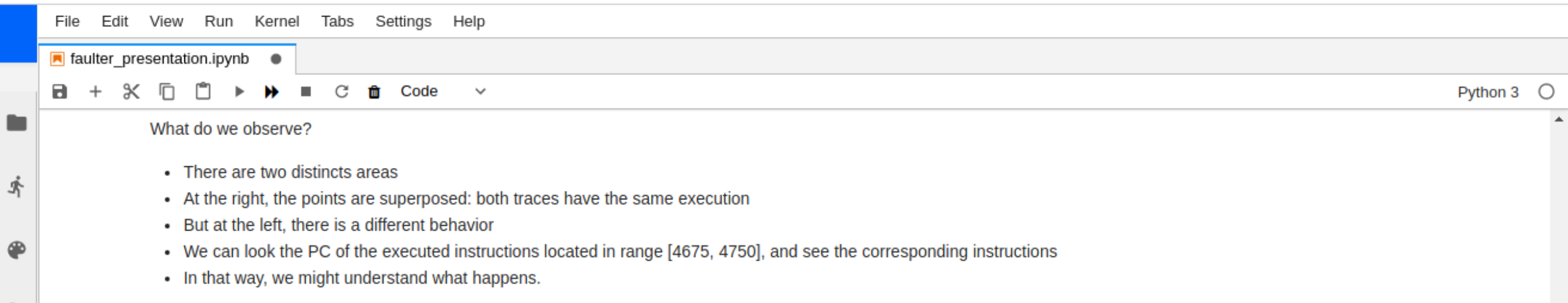
Injecting 1 faults (estimation). Please wait...

Processing 100.0 %  Elapsed: 00:00:00 Remaining: 00:00:00

PC Register - Overview







The image shows a Jupyter Notebook window with a blue header bar. The menu bar includes File, Edit, View, Run, Kernel, Tabs, Settings, and Help. The active tab is 'faulter\_presentation.ipynb'. The toolbar contains icons for saving, adding, deleting, and running code. The text 'Python 3' is visible on the right. The notebook content consists of a text cell asking 'What do we observe?' followed by a bulleted list.

File Edit View Run Kernel Tabs Settings Help

faulter\_presentation.ipynb

Python 3

What do we observe?

- There are two distincts areas
- At the right, the points are superposed: both traces have the same execution
- But at the left, there is a different behavior
- We can look the PC of the executed instructions located in range [4675, 4750], and see the corresponding instructions
- In that way, we might understand what happens.

```
[59]: slideshow('Introduction',  
              './img/presentation_faulter',  
              ['Diapositive%d.PNG' % hit for hit in [49, 50, 51, 52, 54, 55]],  
              slide_width)
```

### ▼ Introduction

Slide 0

Slide 1

Slide 2

Slide 3

Slide 4

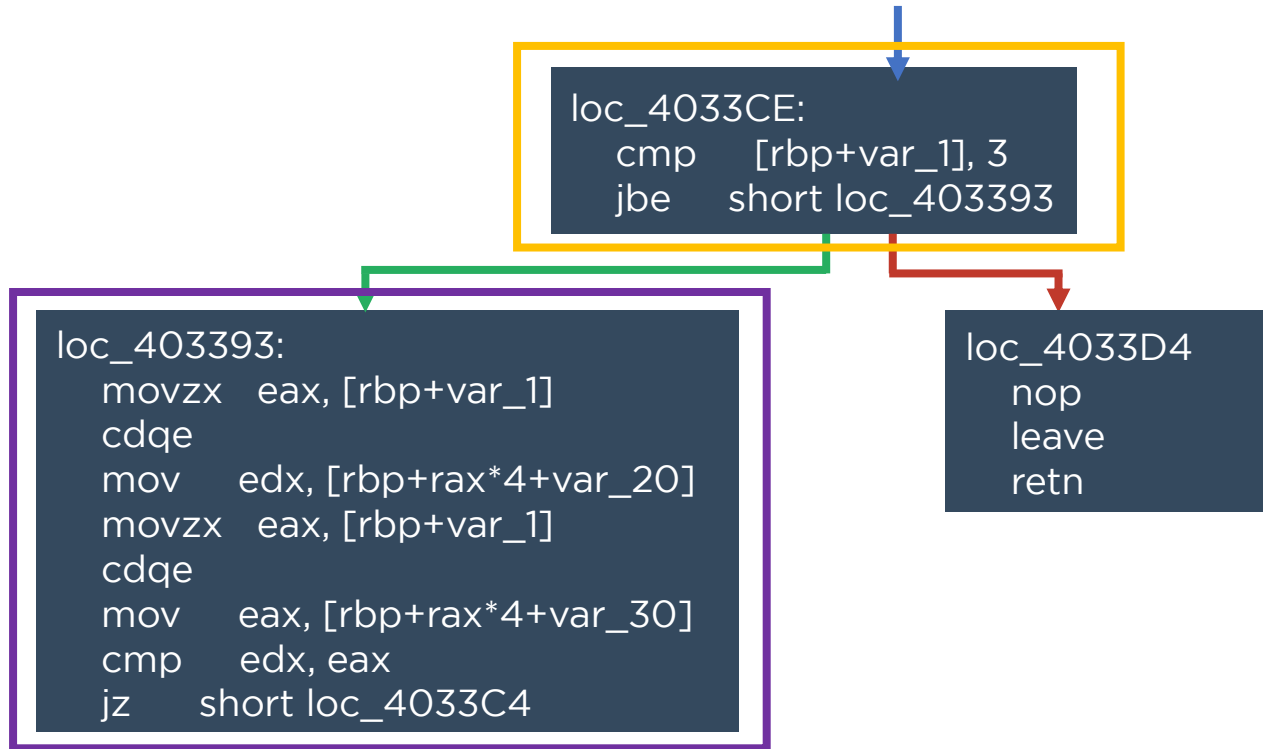
Slide 5

loc\_4033CE:  
cmp [rbp+var\_1], 3  
jbe short loc\_403393

We start the output  
analysis

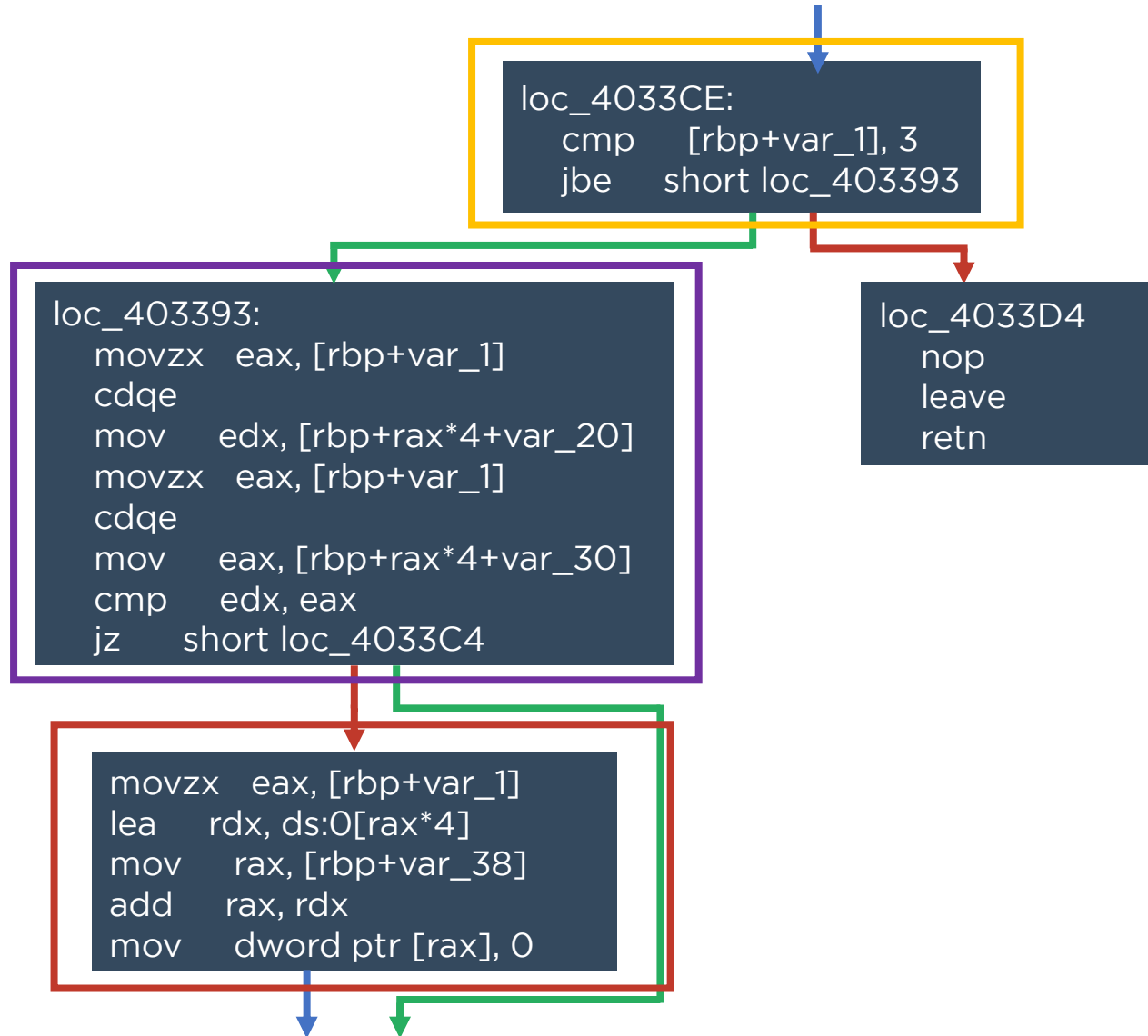


## We start the output analysis



The output was computed twice. Its consistency is checked by block of four bytes.

## We start the output analysis



The output was computed twice. Its consistency is checked by block of our bytes.

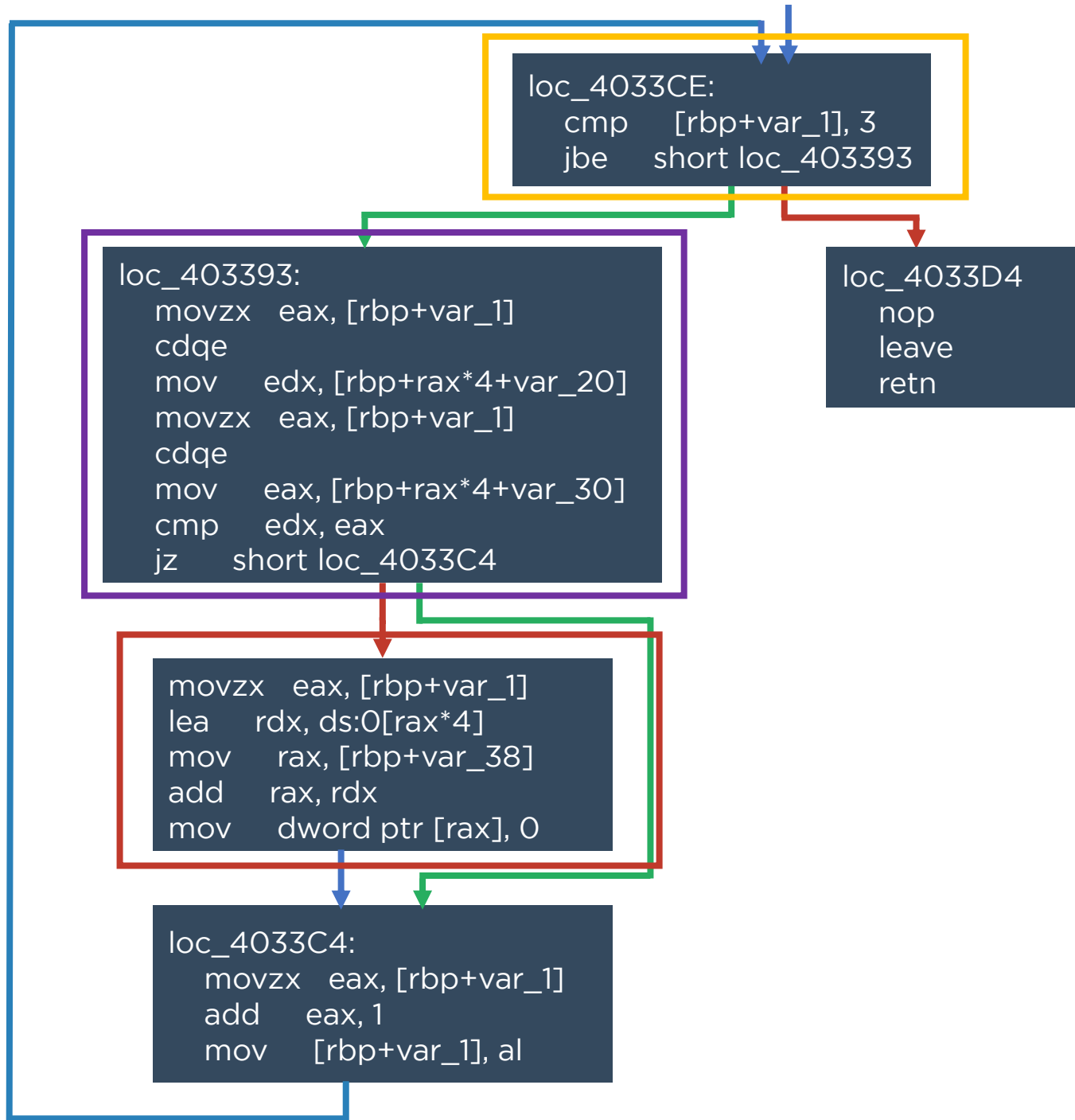
In case of a failure, the four-byte block is set to zero.

We start the output analysis

The output was computed twice. Its consistency is checked.

In case of a failure, the four-byte block is set to zero.

These operation are done 4 times to analyze all the output.



We start the output analysis

The output was computed twice. Its consistency is checked.

In case of a failure, the four-byte block is set to zero.

These operation are done 4 times to analyze all the output.

```
loc_4033CE:
    cmp    [rbp+var_1], 3
    jbe    short loc_403393
```

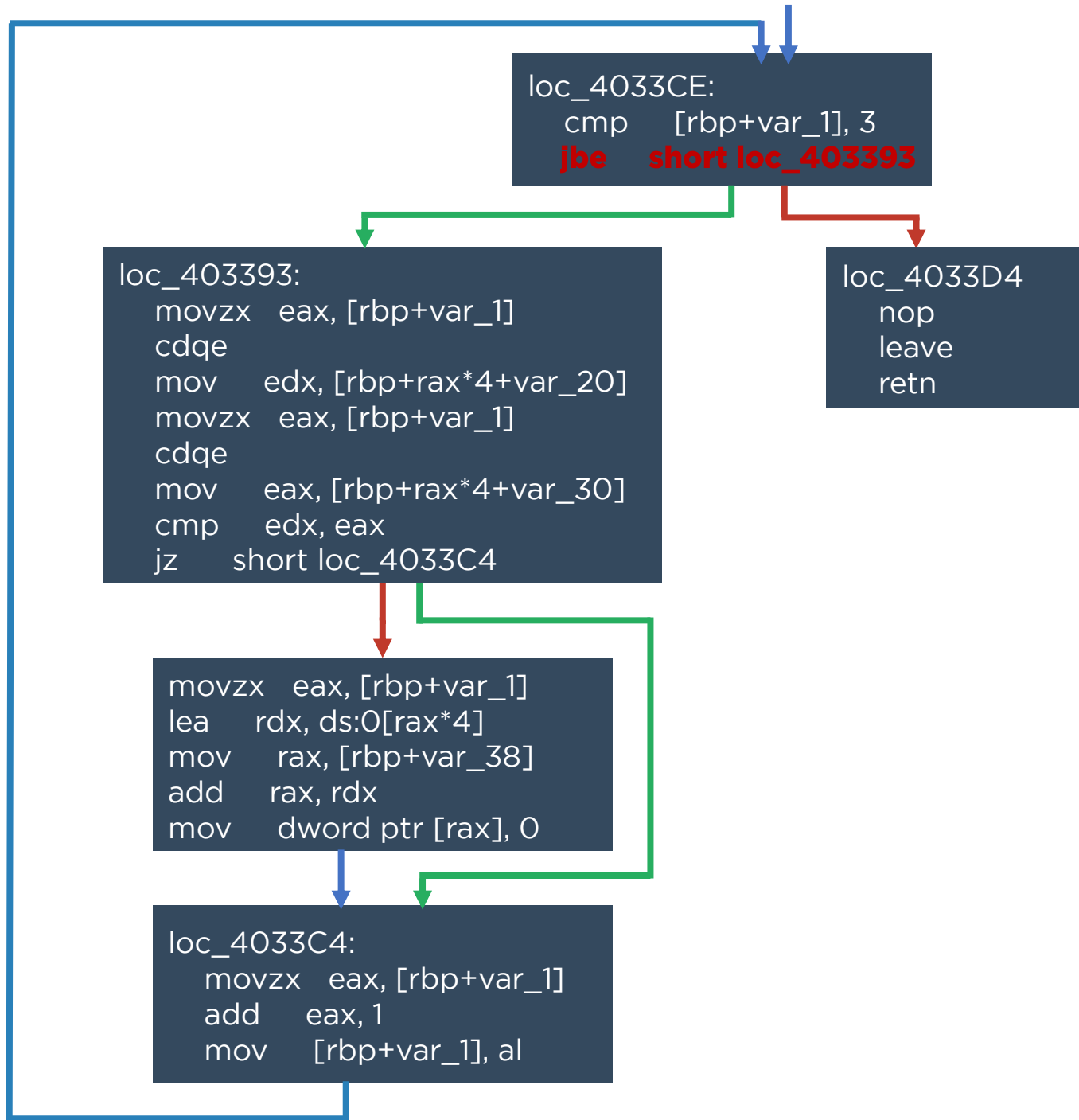
```
loc_403393:
    movzx  eax, [rbp+var_1]
    cdqe
    mov    edx, [rbp+rax*4+var_20]
    movzx  eax, [rbp+var_1]
    cdqe
    mov    eax, [rbp+rax*4+var_30]
    cmp    edx, eax
    jz     short loc_4033C4
```

```
loc_4033D4:
    nop
    leave
    retn
```

```
    movzx  eax, [rbp+var_1]
    lea    rdx, ds:0[rax*4]
    mov    rax, [rbp+var_38]
    add    rax, rdx
    mov    dword ptr [rax], 0
```

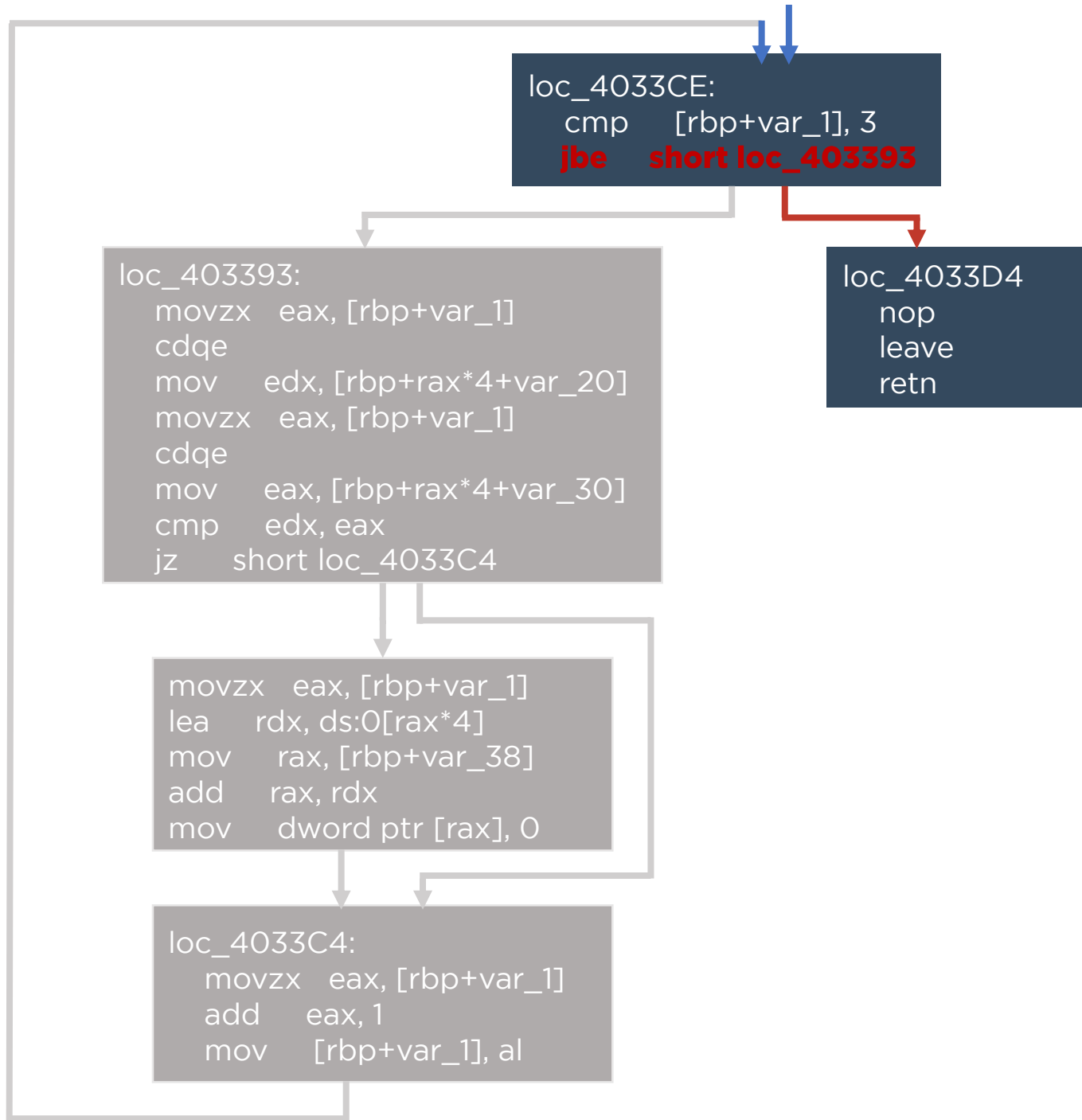
```
loc_4033C4:
    movzx  eax, [rbp+var_1]
    add    eax, 1
    mov    [rbp+var_1], al
```

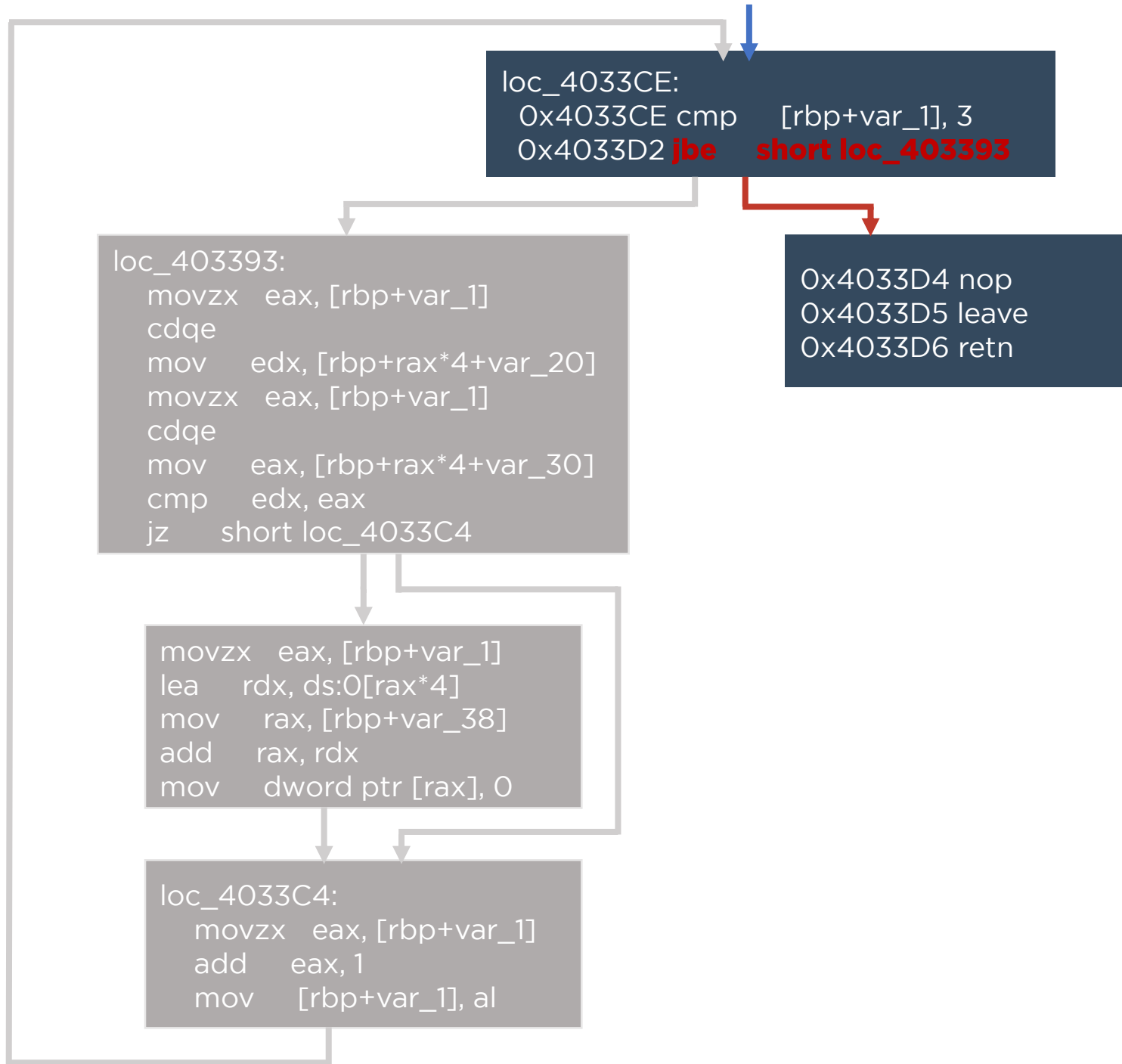
**HOW TO  
DEFEAT THIS  
PROTECTION?**



**We can disturb the control flow:**

- **We jump from jbe instruction directly to 0x4033D4**
- **We exit the function**
- **No check is done**
- **No byte is set to zero**





## 2<sup>nd</sup> fault paramer:

- **Register: Program Counter**
- **When: instruction located at 0x4033D2 is reached**
- **Fault model: Add user value + 2 to jump directly to 0x4033D4**

## 2nd FAULT INJECTION

We are going to inject two faults:

- the first fault aims to set the output to zero
- the second fault targets the consistency verification.

We remove the fault&trace option.

```
[ ]: faulter.tracer = None
```

We will attack only the rax register, and target only the PC that enabled to have an output set to zero.

```
[ ]: pc_to_attack = []
for hit in parser.get_fault_dico()["OUTPUT"]["00000000000000000000000000000000"]["faults"]:
    pc_to_attack.append(int(hit[0]["pc"], 16))
```

We configure the first fault.

```
[ ]: fault_1.register = "rax"
fault_1.filt_pre.in_pc.pts = pc_to_attack[0:32]
fault_1.trig_and_acts = [trig_and_act_start_fault_injection,
                        trig_and_act_stop_fault_injection]
```

We change the directory.

```
[ ]: faulter.directory_out = "faulter_outputs_two_fault"
```

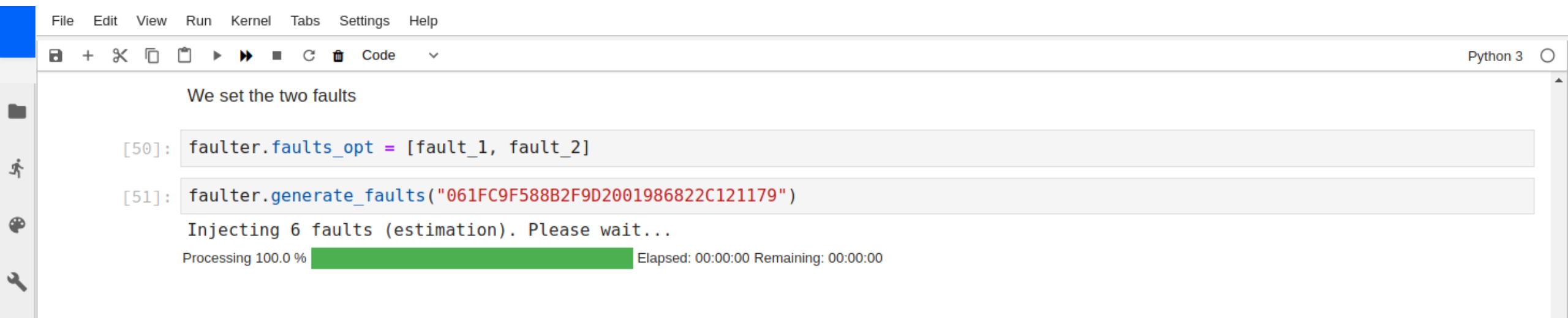
We configure the second fault:

- we want to modify the PC register
- when its value is 0x4033D2, we add 2 to it to skip an instruction and jump directly to 0x4033D4

```
[ ]: fault_2 = FaultCfgMono(bin_hand)
```

```
[ ]: fault_2.filt_pre.in_pc.pts = [0x4033D2]
fault_2.register = "pc"
fault_2.param = 2
fault_2.model = const.FM_ADD_USER_VAL
```





We analyse the new log.

```
[77]: logfile = glob.glob("./s/*.bin" % faulter.directory_out)[0]
      parser = FaulterLogParser(logfile, output_parser = bin_hand.process_output)
      parser.parse_file()
      # Print faulty outputs with 4 modified bytes
      parser.get_faulty_output([4])
```

```
[77]: [['14ed01ea7ce2a551c9791ae85c7cecf4', '6AED01EA7CE2A570C979D3E85CD2ECF4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '4EED01EA7CE2A586C979FAE85C90ECF4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '27ED01EA7CE2A5AEC97950E85C40ECF4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '142301EA11E2A551C9791AC25C7C67F4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '142B01EA7FE2A551C9791A575C7CD9F4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '147101EAC2E2A551C9791AAA5C7C90F4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '145E01EA6CE2A551C9791A5C5C7C53F4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '14ED09EA7C97A551CE791AE85C7CEC98'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '14ED2CEA7CA2A55163791AE85C7CECAC'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '14ED4AEA7C78A5514E791AE85C7CEC30'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '14ED2BEA7C35A55146791AE85C7CEC59'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '14ED01057CE2D551C9F61AE8F77CECF4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '14ED01E17CE2EF51C91B1AE80C7CECF4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '14ED01E47CE22451C99B1AE8447CECF4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '14ED01E87CE24B51C9E51AE8E57CECF4'],
      ['14ed01ea7ce2a551c9791ae85c7cecf4', '14ED01847CE22551C94D1AE80C7CECF4']]
```

This time, we have faulty outputs with 4 modified bytes.

## Key recovery

[53]: `from estoolkit2.fault_attack.aes import DfaAesPiret`

```
# We perform the DFA of Piret
# We need output with 4 faulty bytes
lst = parser.get_faulty_output([4])
dfa = DfaAesPiret(lst)
dfa.run()
```

```
found_key = dfa.get_found_key()
dfa.summary()
```

### DFA summary:

| Words       | 0    | 1    | 2    | 3    | 4    | 5    | 6    | 7    | 8    | 9    | 10   | 11   | 12   | 13   | 14   | 15   |
|-------------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|
| Found bytes | 0x97 | 0x78 | 0x13 | 0xD5 | 0x5A | 0x7B | 0x8A | 0xE1 | 0x55 | 0x8C | 0x31 | 0x3C | 0xA0 | 0xBA | 0x7E | 0x6D |
|             |      |      |      |      |      |      |      |      |      |      |      |      |      |      |      |      |

We found the round key 0x97 0x78 0x13 0xD5 0x5A 0x7B 0x8A 0xE1 0x55 0x8C 0x31 0x3C 0xA0 0xBA 0x7E 0x6D. However, we still need to recover the master key, and then check it.

To recover the master key, we need:

- the extracted round key,
- its round number,

[54]: `from essva.tools.tools import inv_key_schedule_aes`  
`master_key = inv_key_schedule_aes(round_key=found_key, round_nb=10, verbose = True)`

### Computed AES Master Key:

| Words               | 0    | 1    | 2    | 3    | 4    | 5    | 6    | 7    | 8    | 9    | 10   | 11   | 12   | 13   | 14   | 15   |
|---------------------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|
| Round key 10        | 0x97 | 0x78 | 0x13 | 0xD5 | 0x5A | 0x7B | 0x8A | 0xE1 | 0x55 | 0x8C | 0x31 | 0x3C | 0xA0 | 0xBA | 0x7E | 0x6D |
| Computed Master Key | 0x 1 | 0x12 | 0x23 | 0x34 | 0x45 | 0x56 | 0x67 | 0x78 | 0x89 | 0x9A | 0xAB | 0xBC | 0xCD | 0xDE | 0xEF | 0xFF |

We found a key! Let's test it.

```
[55]: message = "096F90870768821A84BD665C2F7DDD4A"
      res = bin_hand.check_key(key=master_key, message=message)
```

### Key matching: SUCCESS

| Words     | 0    | 1    | 2    | 3    | 4    | 5    | 6    | 7    | 8    | 9    | 10   | 11   | 12   | 13   | 14   | 15   |
|-----------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|
| Key bytes | 0x01 | 0x12 | 0x23 | 0x34 | 0x45 | 0x56 | 0x67 | 0x78 | 0x89 | 0x9A | 0xAB | 0xBC | 0xCD | 0xDE | 0xEF | 0xFF |

| Words               | 0    | 1    | 2    | 3    | 4    | 5    | 6    | 7    | 8    | 9    | 10   | 11   | 12   | 13   | 14   | 15   |
|---------------------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|
| Plaintext           | 0x09 | 0x6F | 0x90 | 0x87 | 0x07 | 0x68 | 0x82 | 0x1A | 0x84 | 0xBD | 0x66 | 0x5C | 0x2F | 0x7D | 0xDD | 0x4A |
| Binary ciphertext   | 0x39 | 0x5A | 0xF9 | 0x24 | 0xAC | 0xB5 | 0x2E | 0xB5 | 0xA2 | 0xDC | 0x8F | 0xAF | 0xFC | 0x41 | 0xDC | 0x58 |
| Expected ciphertext | 0x39 | 0x5A | 0xF9 | 0x24 | 0xAC | 0xB5 | 0x2E | 0xB5 | 0xA2 | 0xDC | 0x8F | 0xAF | 0xFC | 0x41 | 0xDC | 0x58 |

It matches! The AES key is 0x01 0x12 0x23 0x34 0x45 0x56 0x67 0x78 0x89 0x9A 0xAB 0xBC 0xCD 0xDE 0xEF 0xFF.

# Conclusion

```
[56]: slideshow('Conclusion',  
             './img/presentation_faulter',  
             ['Diapositive%d.PNG' % hit for hit in [61,62,63 ]],  
             slide_width)
```

## ▼ Conclusion

Slide 0

Slide 1

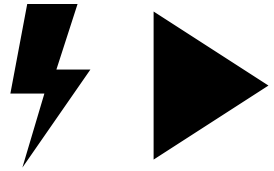
Slide 2

# AGENDA

- Introduction
- How to fault a White-Box
- Requirements to fault a White-Box
- Double Fault injection on an AES White-Box
- **Conclusion**

# CONCLUSION

Faulting a White-Box must be focused on the binary.



Dynamic fault injection is a prerequisite



Accurate multiple faults can be injected



Security mechanisms can be defeated

But

Combinatorial complexity

$$\frac{nb\_ins \times nb\_fault\_model \times nb\_target \times nb\_input \times nb\_area}{nb\_input \times nb\_area}$$

# CONCLUSION

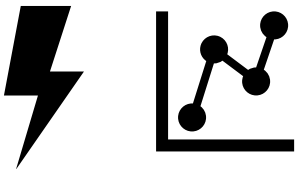
Strategies to defeat these issues:



Pattern detector to fault only interesting area



Focus faulting on specific Register / Program Counter or instructions



Fault & trace to understand the effect of a fault or downgrade security

In that way, it is possible to execute successful multi-fault attacks, in reasonable amount of time.