Secure Protocols in a Hostile World

for CHES 2015

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Why this talk?

10 Things You Should Know About Computer Security

5: Cryptography is a Solved Problem

Cryptography: The strongest link in the chain^{*}

but not to others. Unfortunately, people concentrate too much on the cryptography of a system – which is the equivalent of strengthening the strongest link in a chain.





Unsolved





Why does this matter?



Joto fail; // Apple SSL bug test site

This site will help you determine whether your computer is vulnerable to **#gotc**

Tracking the FREAK Attack



LOGJAM ATTACK (CVE-2015-4000)

TLS Vulnerability

TOP SECRET STRAP1

Response to improving security

- For the past decade, NSA has lead an aggressive, multi-pronged effort to break widely used Internet encryption technologies
- Cryptanalytic capabilities are now coming on line
- Vast amounts of encrypted Internet data which have up till now been discarded are now exploitable
- Major new processing systems, SIGDEV efforts and tasking must be put in place to capitalize on this opportunity

PTD "We penetrate targets' defences."



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Highly Efficient GF(2^8) Inversion Circuit Based on Redundant GF Arithmetic and Its Application to AES Design
Rei Ueno (Tohoku University); Naofumi Homma (Tohoku University); Yukihiro Sugawara (Tohoku University); Yasuyuki Nogami (Okayama
University); Takafumi Aoki (Tohoku University)

2. Robust Profiling for DPA-Style Attacks

Carolyn Whitnall, Elisabeth Oswald (University of Bristol)

3. SoC it to EM: electromagnetic side-channel attacks on a complex system-on-chip

Jake Longo (University of Bristol); Elke De Mulder (Cryptography Research Inc.); Dan Page (University of Bristol); Michael Tunstall (Cryptography Research Inc.)

4. TriviA: A Fast and Secure Authenticated Encryption Scheme

Avik Chakraborti (Indian Statistical Institute Kolkata); Anupam Chattopadhyay (School of Computer Engineering, NTU Singapore); Muhammad Hassan (RWTH Aachen University); Mridul Nandi (Indian Statistical Institute Kolkata)

5. Stealing Keys from PCs using a Radio: Cheap Electromagnetic Attacks on Windowed Exponentiation

Daniel Genkin (Technion and Tel-Aviv University); Lev Pachmanov (Tel-Aviv University); Itamar Pipman (Tel-Aviv University); Eran Tromer (Tel-Aviv University)

6. Efficient Ring-LWE Encryption on 8-bit AVR Processors

Zhe Liu (University of Luxembourg); Hwajeong Seo (Pusan National University); Sujoy Sinha Roy (K.U. Leuven); Johann Großschädl (University of Luxembourg); Howon Kim (Pusan National University); Ingrid Verbauwhede (K.U. Leuven)

This talk

- We know how to build strong cryptosystems
 - Our research focuses on building <u>stronger</u> crypto systems!
- And yet we continue to deploy weak ones
 - Worse, we're largely stuck with weak ones
 - What's going on here?
 - Main case studies: SSL/TLS, IPSEC

Case study I: SSL/TLS

SSL/TLS

- Most important security protocol on the Internet
 - Allows secure connections between clients & servers
 - Current version: TLS 1.2
 - (But browsers still support SSL 3,TLS 1.0/1.1) plus 1.3 coming soon!
 - Not just web browsing!



A brief history

- SSLvI born at Netscape. Never released. (~1994)
- SSLv2 released one year later
- SSLv3 (1996)
- TLS I.0 (1998)
 - Still widely deployed
- TLS I.I (2006)
- TLS I.2 (2008)



How secure is TLS?

- Many active attacks and implementation vulnerabilities
 - Heartbleed, Lucky I 3, FREAK, CRIME, BEAST, RC4





As tomorrow is April 1, today marks the last day of useful e-commerce before SSL breaks again on Thursday. Hope you made the most of it.

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In practice: most of these require substantial resources and <u>can't be deployed at scale</u>





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How secure is TLS?

But not all attacks...



What's wrong with TLS?

Quite a bit

- Many problems result from TLS's use of "<u>pre-historic cryptography</u>" (- Eric Rescorla)
 - CBC with Mac-then-Encrypt, bad use of IVs
 - RSA-PKCS#IvI.5 encryption padding
 - RC4
 - DH parameter generation
 - Horrifying backwards compatibility requirements

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• RC4

Many of these flaws were 'known' at design time, but exploited by researchers only afterwards.

MAC-then-pad-then-Encrypt

- TLS MACs the record, then pads (in CBC), then enciphers
 - Obvious problem: padding oracles
 - Countermeasure(s):
 - I. Do not distinguish padding/MAC failure
 - 2. "Constant-time" decryption

Unlucky for you: UK crypto-duo 'crack' HTTPS in Lucky 13 attack **OpenSSL patch to protect against TLS decryption boffinry** By John Leyden • Get more from this author Posted in Security, 4th February 2013 16:58 GMT

BEAST

- Serious bug in TLS 1.0
- Allows complete decryption of CBC ciphertexts
- Use of predictable Initialization Vector (CBC residue bug)
 - Known since 2002, attack described by Bard in 2005 (Bard was advised to focus on more interesting problems.)
 - Nobody cared or noticed until someone implemented it

Solution in practice: RC4

:-(

(When RC4 is your solution, you need a better problem)

Compression (CRIME)

- Can't really blame the TLS designers for including it...
 - Blame cryptographers for not <u>noticing it's still in use?</u>
 - Blame cryptographers for pretending it would go away.
- We need a model for compression+encryption
 - Clearly this can't be semantically secure
 - But how much weaker? Can we quantify?

Protocol Design

Example: Negotiation

Each TLS handshake begins with a cipher suite negotiation that determines which key agreement protocol (etc.) will be used.



Ciphersuite Negotiation

I support: RSA, DHE, ECDHE, RSA_EXPORT

Negotiate

l choose: ECDHE

Ciphersuite Negotiation

I support: RSA, DHE, ECDHE, RSA_EXPORT







M

I support: RSA, DHE, ECDHE, RSA_EXPORT





I support: RSA, DHE, ECDHE, RSA EXPORT

I choose: RSA_EXPORT












Is it sufficient for the client to support only "strong" ciphersuites, even if the server supports weak ones?



Is it sufficient for the client to support only "strong" ciphersuites, even if the server supports weak ones?

- Let A be the set of KA protocols supported by Client
 Let B be the set of KA protocols supported by Server
- If each KA protocol in $A \cap B$ is a secure KA protocol, is the TLS handshake secure?

TLS for cryptographers

- In <u>CRYPTO 2012 (!)</u> we saw the first paper to successfully analyze TLS-DHE [Jager et al.]
- In CRYPTO 2013 a random-oracle analysis of the TLS-RSA handshake [Krawczyk et al.]
- In CRYPTO 2014 an automated analysis of the full handshake, under a new security model [Bhargavan et al.]

TLS for cryptographers

We do not model ciphersuite negotiation/renegotiation, nor session resumption.

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Theorem

Bhargavan et al. theorem statement:
 Let A be the set of KA protocols supported by Client
 Let B be the set of KA protocols supported by Server

If each KA protocol in $A \cup B$ is a secure KA protocol & there exist PRFs, then the TLS handshake is a secure KA protocol.

Theorem

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If each KA protocol in $A \cup B$ is a secure KA protocol & there exist PRFs, then the TLS handshake is a secure KA protocol.

TLS design/deployment assumes this would be $A \cap B$!

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I support: RSA, DHE, DHE_EXPORT, RSA_EXPORT, ECDHE

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I support: RSA, DHE, DHE_EXPORT, RSA_EXPORT, ECDHE



support:

FREAK [Bhargavan et al.]: Due to a bug in SecureTransport, OpenSSL, SChannel, client accepts export-grade RSA key







@matthew_d_green I am still amazed how three *independent* TLS implementations have the exact same bug.



Solution: Fix implementations

Patch OpenSSL, SecureTransport, SChannel so they will <u>recognize</u> an RSA export key exchange message, barf

(patches rolled out March 2015)

RSA DHE

Apple issues security patches to protect devices from the FREAK bug

by Mariella Moon | @mariella_moon | 22 days ago

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I support:





LogJam [Adrian et al.]: Due to a bug in *the TLS* **protocol** client accepts export-grade DHE key



TLS design/deployment assumptions were wrong, and we knew this for years but failed to properly communicate to the community. TLS design/deployment assumptions were wrong, and we knew this for years but failed to properly communicate to the community.

The community made terrible assumptions and didn't ask us what we thought of them. Then they got mired in backwards compatibility issues and <u>only responded to</u> <u>attacks.</u>

(Joint work: Adrian, Bhargavan, Durumeric, Gaudry, Green, Halderman, Heninger, Springall, Thomé, Valenta, VanderSloot, Wustrow, Zanella-Beguelin, Zimmermann) *to appear* 'CCS 2015

- To exploit the downgrade attack, requires solving a 512-bit DL in real time
 - Initially this seems challenging, but NFS algorithm can be heavily optimized for pre-computation using only prime (p)
 - "Oversieving" increases cost of sieving and storage, but reduces cost of linear algebra step & final "descent"

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Example 3: Negotiation Short term (hack) solution:

Fix OpenSSL, SecureTransport, SChannel so they refuse DHE keys <768 bits

patched in NSS, SChannel, BoringSSL, LibreSSL, SecureTransport

(Took months to accomplish this, since it breaks $\sim 1\%$ of the Internet to make this fix)

RSA_EXPORT, ECDHE

Server has a weak, ephemeral Diffie-Hellman public key

ERR_SSL_WEAK_SERVER_EPHEMERAL_DH_KEY

Hide details

This error can occur when connecting to a secure (HTTPS) server. It means that the server is trying to set up a secure connection but, due to a disastrous misconfiguration, the connection wouldn't be secure at all!

In this case, the server needs to be fixed. Google Chrome won't use insecure connections in order to protect your privacy.

Learn more about this problem.

Long(er) term solutions:

Eliminate 1024-bit DHE (but Java).

Stop using common DHE primes.

Use EU-CMA signatures to validate the protocol transcript. Then you can achieve the $A \cap B$ security the TLS designers originally set out to achieve.

(TLS 1.3 adds such a message, provisionally.)

This picture again

• What's going on here?



This picture again

• What's going on here?



This picture again

• What's going on here?



What is LONGHAUL?

(TS//SI/REL) The LONGHAUL system provides the Extended NSA Enterprise with an end-to-end attack orchestration and key recovery service for Data Network Cipher (DNC) and Data Network Session Cipher (DNSC) traffic. LONGHAUL is extensible to allow for the addition of other Digital Network Intelligence cipher types.

Hypothesis I: LONGHAUL is a database of stolen RSA secret keys

- This works well, but it's boring
- Easy to solve: switch to PFS cipher suites (DHE/ECDHE)

TOP SECRET//COMINT//REL TO USA, FVEY

RSA Exploitation Steps

Is it the key exchange RSA? (server hello)

- If so, is the modulus match a known private key? (server certificate)
 - If so, is there 2-sided collect?
 - \geq If so, do we have:
 - Client Hello
 - Server Hello
 - Client Key Exchange

DECRYPTION!



TOP SECRET//COMINT//REL USA, AUS, CAN, GBR, NZL

Happy Dance!!





Problem

- LONGHAUL also purports to decrypt IPSec/IKE
 - IKE does not use RSA
 - It uses Diffie-Hellman for each connection.



TOP SECRET//COMINT//REL USA, AUS, CAN, GBR, NZL

Type 1: IPSec

- IPSec: IP Security
- Complete paired IKE
 - Common UDP ports: 500 and 4500
- Pre-Shared Key (PSK)
 - Router configuration (good source for PSKs)
- Encrypted Payload (ESP or AH)
 - Next Protocol 50 or 51

What is LONGHAUL?



The breakthrough was enormous, says the former official, and soon afterward the agency pulled the shade down tight on the project, even within the intelligence community and Congress. "Only the chairman and vice chairman and the two staff directors of each intelligence committee were told about it," he says. The reason? "They were thinking that this computing breakthrough was going to give them the ability to crack current public encryption."

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Hypothesis 2: The NSA is breaking 1024-bit DHE

- This sounds completely insane
- Maybe it's not

Breaking DHE at scale

- Breaking DHE == solving the Discrete Logarithm problem
 - In theory, this is too expensive for keys >=768 bits
 - However there is a wrinkle...



IT'S NOT AS HARD AS IT LOOKS

Breaking DHE at scale

- A large percentage of Apache/Java/ISS servers use fixed, hardcoded parameters for DHE
 - IPSec/IKE is even worse: nearly 50% of servers will choose Oakley groups 1 and 2 (768/1024) - generated in 1998
 - NFS is heavily optimized for pre-computation using only the primes
 - With specific pre-computation (\$10s-100s of Million/1 year?) an attacker might be able to break 30-50% of DHE connections with academic levels of computing
 - Approximately 30 core days for final descent



IT'S NOT AS HARD AS IT LOOKS

How do we fix this?

- Eliminate 1024-bit DH
 - This is challenging in TLS, since many machines (Java 7) crash on longer parameter lengths
 - D. Gillmor, new extension to negotiate FF-DHE
- Eliminate DHE altogether
 - Move to ECDHE, which is currently not 100% supported
 - Downgrade to RSA (!)
- Eliminate common primes

Why aren't we fixing this?

Why aren't we fixing this?





Large chunks of the internet have rotted with age. They aren't updated or updateable. Wtf do we do?



Conclusion

- Cryptography is challenging! (duh)
 - We fail to push best practices down to the engineering community
 - They fail to pull best practices from the literature, even years after vulnerabilities are known
 - Cryptosystems continue to become more complex and vulnerable
 - This process is not really tolerable anymore

